

MTX VIDEO MEMORY MAPPING.

1.1

Terms used:

VRAM - Video ram.

VDP - Video Display Processor.

CPU - Central Processing Unit.

MSB - Most Significant Bit

LSB - Least Significant Bit

In using MSB and LSB, I have made the assumption that MSB is the left-most bit of any byte, and LSB is the right-most bit. In my notation MSB = Bit 7. LSB - Bit 0.

MTX Vram memory architecture.

1.2

VRAM on the MTX is managed by the VDP chip, which contains its own auto-incrementing address pointer.

The VRAM is independent from the Z30 processor ram as can be seen from diagram 1.0 above.

It can be seen that there is no apparent direct memory mapping to the video screen which is held in VRAM, however the internal architecture of the VDP chip is such that you can perform full memory mapping of various types through VDP ports 1 and 2.

Port 2 is used for address transfers.

Port 1 is used for data transfers

All addressing throughout VRAM is 14 bit. Address set ups require a two byte transfer with two bits left over. To set up an addressing point the low byte of the address is sent through port 2 first, then the high byte is sent of which bits 0 to 5 are part of the address, the mode (described in section 1.3), being indicated by bits 6 and 7.

MODE.

1.3

The truth table below shows the bit set up required to direct the VD chip to either select 'Write data to VRAM or 'Read data from VRAM'.

BIT	6	:	7	

	1	:	0	'Write data to VRAM'
	0	:	0	'Read data from VRAM'

After the address set up has been made, data bytes can either be input or output along port 1, and because the VRAM is managed by an auto-incrementing register, sequential transfers of data bytes can be performed without having to re-set the address pointer on the VDP chip.

If you wish to perform alternate input and output of data bytes to VRAM you must re-set the addressing mode as appropriate.

Address set ups and data transfers require a certain minimum amount of time between sequential processes. This is 11 micro-seconds between sequential address set-ups, and 8 micro-seconds between sequential data transfers.

If you are using BASIC there will never be any problem, but if you are using machine code this is a point to be aware of.

Listed in the next two sections are routines which will perform the job of "PEEKING" and "POKEING" to the video screen.

Section 1.4 describes the BASIC routines. Section 1.5 and 1.6 describe the assembler code.

1.4

Both of these sections assume that you are using 'TEXT MODE' ie VS 5, which is the normal default BASIC VRAM screen set-up.

For BASIC VRAM memory map see Appendix A..

```
100 REM THESE ROUTINES USE TWO VARIABLES
      VADDRESS - USED TO SET UP VRAM ADDRESS
      VDATA - USED TO RECEIVE OR SEND SCREEN DATA
```

```
110 LET VADDRESS=7*1024: LET VDATA=42
```

(Set up VRAM address pointer to 7K and data byte to numerical value for '*')

```
120 GOSUB 1000
```

(POKE byte onto the screen)

```
130 LET ADDRESS=7*1024
```

(Set up VRAM address pointer to 7K)

```
140 GOSUB 2000 : PRINT VDATA
```

(PEEK byte on screen in location pointed by VADDRESS, set up in line 130, and print result)

```

1000 REM POKE DATA BYTE HELD IN VDATA UNTO SCREEN POINTED TO BY
      VADDRESS
1010 LET TEMP2=INT(VADDRESS/256): LET TEMP1=VADDRESS - (TEMP2*256)
1020 OUT (2),TEMP1
1030 LET TEMP2=TEMP2 OR 64 : LET TEMP2=TEMP2 AND 127
1040 OUT (2),TEMP2
1050 OUT (1),VDATA
1060 RETURN

```

```

2000 REM PEEK DATA BYTE ON SCREEN POINTED TO BY VADDRESS, AND
      RETURN RESULT IN VDATA
2010 LET TEMP2 = INT(VADDRESS/256):LET TEMP1=VADDRESS-(TEMP2*256)
1020 OUT (2),TEMP1
1030 LET TEMP2=TEMP2 AND 63
1040 OUT (2),TEMP2
1050 LET VDATA=INP(1)
1060 RETURN

```

1.5

All of these routines use either registers DE or C. DE is always left unchanged as are all other registers, except C which will change dependant on the values you are reading into it when you call VDINPT.

Output a byte.

```

LD      DE,7168      ;TOP OF TEXT SCREEN
CALL    VSETOT      ;SET UP VRAM ADDRESS POINTER FOR DATA
                          ;OUTPUT
LD      C,42        ;NUMERICAL VALUE OF '*'
CALL    VDOUTP      ;OUTPUT BYTE TO SCREEN

```

Input a byte.

```

LD      DE,7168      ;TOP OF TEXT SCREEN
CALL    VSETRD      ;SET UP VRAM ADDRESS POINTER FOR DATA
                          ;INPUT
CALL    VDINPT      ;READ BYTE FROM SCREEN - BYTE RETURNED
                          ;N C

```

VDP I/O routines.

1.6

```

;VSETOT-SET UP VRAM ADDRESS POINTER FOR DATA OUTPUT DEPENDANT
      ON ADDRESS HELD IN DE ON ENTRY

```

```

;
;VSETOT:  PUSH      AF      ;SAVE ACC
          LD        A, E
          OUT      (2),A    ;OUTPUT LOW BYTE ADDRESS
          LD        A,D
          OR      64

```

```

        AND      127      ;SET 'WRITE TO VRAM MODE'
        OUT      (2),A    ;OUTPUT HIGH BYTE ADDRESS
        POP      AF       ;GET OLD ACC
        RET

;
;VSETRD - SET UP VRAM ADDRESS POINTER FOR DATA INPUT DEPENDENT
;          ON ADDRESS HELD IN DE ON ENTRY
;
VSETRD:  PUSH     AF       ;SAVE ACC
        LD      A, E
        OUT     (2),A     ;OUTPUT LOW BYTE ADDRESS
        LD      A, D
        AND     63       ;SET 'READ FROM VRAM MODE'
        OUT     (2),A     ;OUTPUT HIGH BYTE ADDRESS
        POP     AF       ;GET OLD ACC
        RET

;
;VDOUTP-OUTPUT DATA BYTE HELD IN C TO ADDRESS POINTED TO BY AUTO-
;          INCREMENTING REGISTER ON-BOARD THE VDP
;
VDOUTP:  PUSH     AF       ;SAVE ACC
        LD      A, C
        OUT     (1),A     ;OUTPUT DATA BYTE
        POP     AF       ;GET OLD ACC
        RET

;
;VDINPT-INPUT A DATA BYTE AND RETURN VALUE IN C FROM ADDRESS
;          POINTED TO BY THE AUTO INCREMENTING REGISTER ON BOARD
;          THE VDP
VDINPT:  PUSH     AF       ;SAVE ACC
        IN      A,(1)     ;READ DATA BYTE
        LD      C,A      ;PLACE BYTE IN C'
        POP     AF       ;GET OLD ACC
        RET

```

Using the MTX assembler.

1.7

The MTX assembler is a simple to use in line, assembler called from BASIC. Only the machine executable OBJECT code is stored in memory, readable assembler SOURCE code is generated by disassembling the object code and inserting the relevant text and labels, stored in tables below the object code.

The actual location of the code can change as it is stored as a BASIC fine, within BASIC. So as a program is extended and edited the location of the code changes.

Writing in Assembler.

As assembler code is stored in a BASIC line it is first necessary to tell the computer at which BASIC line you wish the code to appear. This is done with the immediate command ASSEM, eg :

ASSEM 10

where 10 is the BASIC line number at which the code will appear.

On pressing the return key the screen will clear and the prompt :

Assemble >

Will appear at the bottom of the screen. You are now in assemble mode !

The MTX will no longer respond to BASIC commands such as LIST or RUN but instead expects one of the assembler commands:

L. - List	T - Top of program
E - Edit	Insert (by default)

Creating a Program.

To start writing your program press the RET key. The word "Insert" will appear at the bottom of the screen.

A few lines higher up a 4 digit hex number representing the memory location to which the next instruction will be assigned, the flashing cursory and the instruction RET appears.

Press the EOL key on the cursor control keypad to get rid of the current instruction and enter your own.

When you have typed and edited your assembler line press the RET key. The line you have just typed in will disappear into memory and the next line will appear ready for another assembler mnemonic.

When you have finished writing your program press CLS on the cursor control keypad followed by RET. The Assemble prompt should re-appear.

At this point you may wish to list your program out on screen. Press T and RET to set the program location pointer to the top of the program followed by L and RET to list out the program.

If your program is more than can be displayed on one screen then after the first screen has been presented a bell will ring and the listing will stop. To continue the listing press any of the keys on the main key board. Should you want to stop the listing then press the BRK key on the cursor control key pad.

When the listing is finished the assembler is again ready to accept another command, ie. the Assemble prompt appears at the bottom of the screen.

In the following program the data stored in the five bytes starting at DATA are transferred to the five bytes starting at COPY.

```
4007      LD HL,DATA
400A      LD DE,COPY
400D      LD BC,5
4010      LDIR
4012      RET
4013 DATA: DB 12,£34,"LOW"
4010 COPY:DS 5
4011      RET
```

Symbols..

```
DATA 4013  COPY      4018
```

Suppose we wanted to modify the program to transfer only the first two bytes. We would change the line at 400D from LD BC,5 to LD BC,2.

To do this type E £400D and RET.

This will put us in edit mode and display the line to be edited.

Use the cursor along the line with the arrow keys until it is over the letter "5", then we press the "2" key to change the number followed by RET to store the line away. The line is now safe in memory and the next line will be presented for editing.

To quit edit mode type CLS followed by RET.

This will leave the assembler ready to accept the next directive.

There is another way of specifying lines in the assembler code without having to use the address of the instruction.

Suppose we wanted to alter the 12 in line 4013 of the previous example. That line has an associated label "DATA", so we can use this instead of the address.

Type E DATA and RET.

This has exactly the same effect as E £4013 and RET.

To leave the assembler press CLS and RET. This will bring you back to BASIC.

Now type LIST or L. in BASIC and you will see your assembler code program appear as a BASIC line. If you already had or now add some lines of BASIC you will find the code line takes its place just like any normal BASIC line would.

NB:- as your BASIC program is added to or edited the code line will move about in memory. As most machine code is location dependent it will become necessary to re-assemble your code.

To do this type:... ASSEM 10

and then carriage return to get into the code line and summon the assembler followed by CLS and RET to leave it again.

This will re-assemble the code for its current location. To get around this problem it is wise to put your code lines as near the start of the BASIC program as possible. eg:

```
10 GOTO 100
100 CODE
4007 LD HL,DATA
400A LD DE,DATA1
etc.
Symbols.
DATA 4013 COPY 4018
20 RETURN
100 REM START OF BASIC PROGRAM
```

Number Representation.

The MTX assembler will work with either decimal or Hexadecimal numbers. The default is decimal, but by prefixing a hex number with a "£" , the number is treated as being hexadecimal. Words are stored according to the Z80 convention, low byte first, then high byte.

Executing Machine code.

There are two ways to execute machine code from BASIC.

The first and simplest is to place the code "in the way" of the BASIC program execution flow. When the code line is encountered by BASIC, control will be passed to the machine code in the code line, control is returned to BASIC by the RET statement. Should you wish to execute the code more than once then it is best to incorporate the code line in a BASIC subroutine which can then be called over and over. This also has the advantage of allowing you to place the code line near the beginning of the BASIC program and avoiding having to re-assemble the code every time the BASIC program is edited.

The second method of executing machine code is by the use of the USR function. Take the previous example and add the following BASIC lines to it..

```
5 GOTO 100
100 LET X=USR(16400)
110 STOP
```

If your machine is an MTX500 then line 100 should be:

```
100 LET X=USR(32784)
```

Now re-assemble the code line to take account of the fact it has been moved up the memory by inserting line 5.

Type ASSEM 10 - Carriage return - CLS -Carriage return

The program can now be "RUN". The number in brackets in the USR function is the address of the machine code, and needs to be re-calculated if the code is moved.

USR returns with the value of the BC register pair, which it.) this case, is assigned to X. BC will be zero after execution of the routine.

Assembler Commands.

The assembler has only 4 different commands, L,E,T and an insert command which works by default. The syntax for E, L and insert is :-

<command letter L,E> space <hexadecimal number> or <label>

Where command letter refers to L, E or in the case of insert no letter.

- 1) Any number used must be within the range of the code line.
- 2) Any label used must be contained within the current code line.

The space between the command letter and its operands is optional, but if the command is used with a label then this may be misunderstood.

eg:- If you wish to edit the line DIOT then using E DIOT is ok but using EDIOT when the label "EDIOT" exists will cause the assembler to insert lines before the label EDIOT.

If no parameter is given with the command then the default is the current program line, ie the line after the fast amended or inserted line. This can be quite useful when switching between edit and insert modes.

The T command has no parameters and is used only to set the program line location pointer to the start of the program.

Pseudo Ops.

Pseudo operations look like Z80 instruction mnemonics but are not. They are used to define the contents of memory and reserve space.

There are only 3 pseudo operations in the MTX assembler, DB,DW,DS:

- 1) DB - Define byte.

This instruction is used to define the contents of a particular byte in memory. Its arguments can be numeric, either decimal or hex, alphanumeric, or the low byte of a label.

Eg:- <label>: DB <dec/hex No.> and/or "<string>" and/or <label>

NB:- Using a label will generate an out of range error, but ignoring this error leaves the low byte of the label in memory.

2) DW - Define word.

DW allocates a value to word ie two bytes. The number is stored low byte first, making it compatible with Z80 word instructions.

Eg:- <label>: DW <dec/hex.no.> and/or <label>

3) DS - Define space.

This reserves a specified amount of space in memory, which may be 0 to 254 bytes.

Eg.- <label>: DS <dec/hex.No.>

Examples of Pseudo ops.

```
DATA:      DB 10,£20, "HIGH"  
           DB "SCORE"  
JMPTAB:    DW START,START1,0,£FOE3  
           DW HIT,WIN  
BUFFER:    DS 50  
           DS £40
```

Inserting comments.

The MTX assembler accepts comments in the usual way, ie prefixed by a ";" and delimited by the end of the line.

If you wish to insert comments after non-executable lines you must use NOP,s.

```
NOP        ;Routine to show comments  
NOP  
ROUTE:     INC A  
           RET
```

Listing, Loading and Saving.

As the assembler code is stored as a BASIC line there is no problem in listing loading and saving, these are all done as for BASIC.

MTX 500 and 512.

The main difference between the MTX 500 and 512 is the way in which RAM pages are configured.

The 500 has 32k bytes of memory starting from 9000 hex and finishing at top of memory at FFFF hex.

The 512 has 64k of memory starting at 4000 hex and finishing at FFFF hex, with an additional page of 16k switched out between 3000 hex and C000 hex. This additional memory is switched in and out automatically and provided the code line does not pass over its boundary the user is unaware of the paging system which is maintained in hardware by a ULA.

It is possible to trick an MTX 512 into acting like a 500. To do this go into PANEL and change the contents of memory location FA7A hex (61422 decimal), to zero, then leave PANEL and type NEW.

This will leave the machine ready to load or have typed a new program as though it were a 500.

There are a number of differences between the MTX assembler and other assemblers.

Among these are the fact that there is no requirement to specify the origin of the object code or provide a long list of assembler directives. As a result the assembler is both easy to use and relatively fast, about 2 or 3 seconds for an 8k program.

The set of Z80 instructions concerned with loading the stack pointer with HL, IX and IY are not directly available from the assembler.

They can be accessed in the following way.,

Normally	Use instead
LD SP, HL	DB £ F9
LD SP,IX	DB £DD,£F9
LD SP,IY	DB £FD,£F9

The VDP Chip.

Introduction.

2. 1.

The VDP chip has four video display modes. Normally only Text mode and Graphics mode 2 are available directly from BASIC, but you can access the other two modes, Graphics mode 1 and Multicolour mode, by creating your own VDP set-ups (See sections 2.2 to 2.5).

Text mode provides a screen which is 24 times 40 character rows in two colours.

Multicolour mode provides a 64 times 48 colour dot display in 15 colours plus transparent and 32 sprites.

Graphics 1 mode provides a 256 times 192 pixel display in 15 colours plus transparent and 32 sprites.

Graphics 2 mode is an enhancement of Graphics 1 mode providing more complex colour and pattern displays and 32 sprites.

The video display consists of 35 planes, numbered from 34 down to 0. Working from the 'back' of the screen to the 'front' these are listed below.

Plane number

34	External video plane	- (Not used at present)
33	Backdrop plane	- (This is where the border lives)
32	Pattern plane	- (This is where the screen you print appears)
31	Sprite plane 31	- (This sprite has the lowest priority and will be hidden by any of the other 31 sprites which may cross it)
30-1	Sprite planes 30 to 1	
0	Sprite plane 0	- (This sprite has the highest; priority and will appear to pass in front of any of the other 31 sprites it may cross)

It is important to note that the sprite planes are not active in Text mode.

Write only registers.

2.2

The VDP has eight write-only registers and one read only status register. The write-only registers control VDP operation and determine the way in which VRAM is allocated. Section 2 is concerned with the operation of these write only registers and does not cover the operation of the read-only register. For brief description of how this works and what it does see Appendix B.

Each of the eight write-only registers can be loaded using two 8-bit data transfers from the CPU. The first byte transferred is the data byte and the second is the control byte. The data byte can be any value between 0 to £FF (Decimal 0 to 255). The format of the control byte to write to each of the write-only registers (numbered 0 to 7), is shown in the table below.

Register Number	Bits 7, 6	Bits 5, 4, 3, 2, 1, 0	Hex byte	Decimal byte
0	1 0	0 0 0 0 0 0	£80	128
1	1 0	0 0 0 0 0 1	£81	129
2	1 0	0 0 0 0 1 0	£82	130
3	1 0	0 0 0 0 1 1	£83	131
4	1 0	0 0 0 1 0 0	£84	132
5	1 0	0 0 0 1 0 1	£85	133
6	1 0	0 0 0 1 1 0	£86	134
7	1 0	0 0 0 1 1 1	£87	135

It can be seen that bits 7 and 6 MUST be set to 1 and 0 respectively. These are the active control bits in this byte and tell the VDP that the previously transmitted data byte is to be directed into one of the write-only registers. Bits 5,4,3 MUST be set to zeroes. The number of the destination register is indicated by bits 2,1,0.

Both data transfers and control bytes directed to the VDP write-only registers must be output via port 2.

Example 1.

Assuming we are in text mode, the foreground and background colour is determined by write-only register 7. To alter the foreground colour to black and the background colour to white we could use either one of the two sections of code listed below.

The way in which VDP register 7 works is described in section 2.3 register 7.

In BASIC.

```
100 COLOUR 1,1 : COLOUR 2,15
```

Which will perform the same job as the assembler code listed below.

```

                LD    DE,£071F          ;Register D = Write only reg
                                                ;destination ie integer X
                                                ; 0 =< X =< 7
                CALL  VOUTRG           ;Register E = Data byte
                                                ;Output data byte to write only register
;
;
;VOUTRO-OUTPUT A DATA BYTE TO A SPECIFIED REGISTER
;
;    REGISTER DESTINATION HELD IN D
;
;    DATA BYTE HELD IN E
;
;    DE MUST HOLD VALID CONTENTS ON ENTRY
;
;    NO REGISTERS AFFECTED ON EXIT
VOUTRG:  PUSH    AF                    ;Save Acc and flags
                LD    A,E              ;Get data byte
                OUT   (2),A            ;Output data byte
                LD    AA                ;Get control byte
                AND   7                 ;Set up correct control bits
                OR    128               ;Set bits 6,5,4,3 to zero
                OUT   (2),A            ;Output control byte
                POP   AF                ;Get old acc and Flags
                RET                     ; Return to calling routine

```

How BASIC initialises the 2.3

When the MTX is switched on and the VDP write-only registers are set-up for the first time, it is necessary to set and reset write-only register 0 to 'wake up' the VDP chip and then to set up all of the other registers correctly.

BASIC sets up the VRAM tables as shown in Appendix A. The data bytes which are used, and the routine to perform this set-up are listed in the example below.

The routine is called VRGINI (VDP Registers Initialisation Routine). It needs no register setup on input, and affects none on output. It uses a table called VRGTAB which has 18 elements, and the routine VOUTRG, which must be present.

```

;VRGINI-INITIALISE VDP WRITE-ONLY REGISTERS DEPENDENT CONTENTS OF
  TABLE VRGTAB
;
; NO REGISTER SET-UP REQUIRED ON ENTRY
;
; NO REGISTERS AFFECTED ON EXIT
;
;
VRGINI:  PUSH      AF          ;Save Acc and flags
         PUSH      DE          ;Save DE register pair
         PUSH      HL          ;Save HL register pair
         LD        B, 18       ;Set loop counter = 18
         LD        HL,VRGTAB   ;Set HL to point to start address of table
                                   ;VRGTAB
VRG1:    LD        E,(HL)      ;Load data byte into E
         INC       HL          ;Move onto next byte in table VRGTAB
         LD        D, (HL)     ;Load control byte into D
         INC       HL ;       ;Move onto next byte in table VRGTAB
         CALL      VOUTRG      ;Output data byte held in to register number
                                   ;held in D
         DJNZ     VRG1        ;Decrement loop counter and If loop counter
                                   ;<> 0 then goto VRG1 else drop through to
                                   ;VRG2
VRG2:    POP       HL          ;Get old HL register pair
         POP       DE          ;Get old DE register pair
         POP       AF          ;Get old Acc and flags
         RET                          ;Return to calling routine
;
VRGTAB:  DB        0,0,0,0,0,0,0,0,0
         DB        0,0,0,0,0,0,0,0,0
;

```

The contents of the eight VDP write-only registers are shown in the table below. Following this table is a description of each register and what each of the control bits actually do.

Registers 0 and 1 contain flags to enable or disable various VDP features and modes. Registers 2 through to 6 contain values which are the start addresses for the various sub-blocks of VRAM (for example where the screen is located).

Register 7 is used to define backdrop and text colours.

Register Number	MSB	7	6	5	4	3	2	1	LSB	0
0	0	0	0	0	0	0	0	M3	EV	
1	4/16K	BLANK	IE	M1	M2	0	0	SIZE	MAG	
2	0	0	0	0	0	NAME TABLE BASE ADDRESS				
3	COLOUR TABLE BASE ADDRESS									
4	0	0	0	0	0	0	PG BASE ADDRESS			
5	0	SPRITE ATTRIBUTE TABLE BASE ADDRESS								
6	0	0	0	0	0	0	SPG BASE ADDRESS			
7	TEXT COLOUR ONE				TEXT COLOUR ZERO/BACKDROP					

Register 0 This contains two VDP option control bits. All other bits are reserved for future use and must be zeroes.

Bit1 M3 Mode Bit 3 (See register 1, bits 3 and 4)

Bit 0 EV External Video Enable/Disable 1 Enables External Video Input 0 Disables

Register 1 This contains seven VDP option control bits. Bit 2 is reserved for future use and must be zero.

Bit 7 4/16k Ram Selection.
1 selects 4108/4116 RAM operation.
0 selects 4027 RAM operation.

Bit 6 This bit must be set to 1 on the MTX. It uses 4116 RAM chips for VRAM.

BLANK enable/ disable
1 enables the active display
0 causes the active display area to blank
When this bit is set to zero, the screen will show the border colour only

Bit 5 IE = Interrupt enable/disable
1 enables VDP interrupt
0 disables VDP interrupt

This subject is covered in depth in section 3

Bit 4 M1 = Mode bit 1

Bit 3 M2 = Mode bit 2

The three mode bits M1, M2 and M3 determine the graphic mode that the VDP is currently in. The necessary bit set up required to activate the different modes is shown in the table below.

M1	M2	M3	
0	0	0	Graphic 1 mode
0	0	1	Graphic 2 mode
0	1	0	Multicolour mode
1	0	0	Text mode

Bit 2 Reserved and must be zero.

Bit 1 SIZE selects Sprite Size
 1 selects size 1 sprites (16 times 16 bits).
 0 selects size 0 sprites (8 times 8 bits).

Bit 0 MAG is the magnification option for sprites.
 1 selects MAG1 sprites (SIZE times 2).
 0 selects MAG0 sprites (SIZE times 1).

Register 2 The contents of this register define the start address of the name table sub block, in other words, the start address of the VRAM display screen.

It is a four bit value and therefore has a range of 0 to 15, which it must not exceed. Bits 7,6,5 and 4 of this register must be zero.

The name table must be located on a 1k boundary, so if we were to load this register with say 7, the VRAM display screen would start at 7K or £1C00 Hex. If we loaded it with 15, the VRAM display screen would start at 15K or £3C00 Hex. In text mode this table is 960 (24 times 40) bytes long.

In all other modes this table is 768 (24 times 32) bytes long.

Register 3 The contents of this register define the start address of the colour table sub block. This holds the colours of the patterns used in graphic modes 1 and 2. It is not used by text or multicolour mode. Text and multicolour mode are covered in sections 2.5 and 2.7 respectively.

Graphic modes 1 and 2 access the colour table in different ways. This means that the set up bits for this register are different dependant on which graphic mode you wish to use.

In graphic one mode the colour table is located on a 64 byte boundary and is 32 bytes long. If we wanted to locate it at 8k (£2000 Hex), we need to divide 8k (3192 Decimal) by 64, and then load the result into register 3. This will be 32.

For a fuller description of how graphic mode 1 works see section 2.6.

In graphic two mode, which is the usual BASIC graphic mode, the colour table is located on an 8k boundary. As the VDP uses 16K of VRAM, this means that the colour table can be located at either address 0K or 8k only.

If you wish to locate the colour table at address 0K then bit 7 of register 3 must be set to 0. If you wish to locate the colour table at address 8k then bit 7 of register 3 must be set to 1.

In both cases bits 6 to 0 must be set to 1's. This set-up is also shown in the table below.

Locate colour table at VRAM address	Load register 3 with data byte
OK - £0000 – 0 Decimal	127 Dec - £7F Hex
BK – £2000 – 8192 Dec	255 Dec - £FF Hex

The colour table is 6144 bytes long in graphic mode.

Graphic 2 mode colour table creation is described in much greater depth in section 2.4.

Register 4 PG BASE ADDRESS is the start address of the pattern generator table. This is where the patterns of the characters for the various modes are held. It is a 3 bit value and therefore has a range of 0 to 7. It must not exceed this value.

Bits 7,6,5,4 and 3 must be set to 0.

In all graphic modes except graphic mode 2, this table is located on a 2K boundary. This gives 8 possible positions in which it can be located in VRAM. The start address is determined by the contents of this register times 2K.

If for example we wished to locate this table at address 8k (£2000 Hex) in VRAM, we would load register 4 with 8k divided by 2K which is 4.

In graphic 2 mode this table is located on an 8k boundary as in the graphic 2 mode for register 3.

Bits 1 and 0 must be 1. If bit 2 is 1 then the pattern generator table will be located at address 8K . If this bit 2 is 0 then the table will be located at 0K.

This set-up is also shown in the table below.

Locate pat gen table at VRAM address	Load register 4 with data byte
OK - £0000 – 0 Decimal	3 Dec - £03 Hex
8K - £2000 – 8192 Dec	7 Dec – £07 Hex

The pat gen table is 6144 bytes long in graphic 2 mode.

Graphic 2 mode pattern generator creation is described in much greater depth in section 2.4.

Register 5 SPRITE ATTRIBUTE TABLE BASE ADDRESS is the start address for a 128 byte block which contains position, colour and shape information, for each of the 32 sprites which it is possible to enable.

The sprites are active in all modes except text mode.

This register value occupies 7 bits and therefore has a range of 0 to 127. Bit 7 must be set to 0. The sprite attribute table is located on a 128 byte boundary. If we wished to locate this table at say address 6K (£1000 Hex - 6144 Dec), we would have to divide our desired decimal start address by 128 (or our desired hex address by £80), and then place the result in register 5. In this case the correct resulting byte is shown in the table below.

	Desired address	Divide by	Result
Dec	6144	128	48
Hex	1800	£80	£30

For a fuller description of the sprite attribute table please see section 4.

Register 6 SPG BASE ADDRESS is the sprite pattern generator base address. It holds a library of sprite shapes which can be 'called up' by one of the control bytes in the sprite attribute table. This means that to change a sprite shape to another existing pattern it is only necessary to change one byte in VRAM.

This table is a maximum of 2048 bytes long, and is divided into 256 blocks of 8 bytes each.

The register value for this table occupies 3 bits and therefore has a range of 0 to 7. It must not exceed this range. Bits 7,6,5,4 and 3 must all be set to 0.

The table is located on a 2K boundary so if we wished to locate this table at say address 14K, all we have to do is divide the desired start address for this table by 2K to obtain the correct byte to load into register 6. In this case 14K divided by 2K = 7.

For a fuller description of the sprite pattern generator table please see section 4.

Register 7 This register is split into two nibbles ie two 4 bit values.
The upper 4 bits 7,6,5 and 4, contain the colour code of the ink colour for characters in text mode.

The lower 4 bits 3,2,1 and 0, contain the colour code of the paper colour for characters in text mode and the backdrop (border) colour in all modes.

Graphics Mode 2 2.4

Graphics 2 mode is the normal BASIC graphics mode. The features which it provides are summarised in the table below.

Screen 768 unique characters (24 rows by 32 columns). 256 (Horizontal) by 192 (Vertical) plottable pixels.

15 colours (See appendix C) plus transparent, in a 3 (Vertical) by 2 (Horizontal) colour matrix for each character.

15 border (Backdrop) colours plus transparent.

Sprites 32 sprites each one of which may have any one of 15 colours including transparent and are plottable at any of the 256 times 192 screen pixel positions.

256 patterns for 8 pixel by 8 pixel sprites. Bits 1 and 0 in VDP write-only register 1 are both 0.

64 patterns for 16 pixel by 16 pixel sprites. Bits

1 and 0 in VDP write only register 1 are 1 and 0 respectively.

The VDP is in graphics 2 mode when mode bits M1 = 0, M2 = 0 and M3 = 1 (see section 2.2 register 1 bits 3 and 4).

When the DP is first initialised into graphics two mode, the VRAM is organised as displayed in the table below.

VRAM sub block	Length in bytes
Pattern generator table	6144
Pattern colour table	6144
Pattern name table (screen)	768
Sprite attribute table	128
Sprite generator table	2048

If you had created your own graphic 2 mode set up then these tables could be located at various addresses. If however, you had set them up at the same addresses as BASIC does, then the VRAM memory map for these tables would look like that below.

Address to decimal

0000	- start of pattern generator table
6143	- end of pattern generator table
6144	- free space (see below)
8191	- end free space. ^ (Illus 2 – VMM)
8192	- start of pattern colour table
14335	- end of pattern colour table
14336	- start of sprite generator table
15359	- end of sprite generator table
15360	- start of pattern name table
16127	- end of pattern name table
16128	- start of sprite attribute table
16255	- end of sprite attribute table

It can be seen that the sprite generator table as set up by BASIC is only 1K bytes long whereas it is normally 2K bytes long. It means that you can only use half of the number of possible sprite patterns that would normally be available because this table overlaps with the display screen.

As it has to be located on a 2K byte boundary and as there is apparently 2K bytes free at address 6144 Decimal, most readers will be wondering why the sprite generator was not located there instead.

The reason for this is that this area is actually reserved for text mode which occupies the 'free space' as shown below.

Address in decimal

6144	- text pattern library
7167	- end text pattern library
7168	- start text name table (text display screen)
7191	- end text name table

Allocating 1K bytes for the sprite generator table as BASIC does is a compromise. The VRAM set up as BASIC creates, it means that to switch from text mode to graphics 2 mode and vice versa it is only necessary to change two VDP write-only registers, a process which takes a little over 20 micro-seconds (20 1/100 thousandths of a second).

This is how you can switch from one mode to the other without affecting the integrity of either screen in any way. The tables are held separately, and when the modes are switched, are left intact because they have clearly defined separate boundaries within VRAM.

In graphics 2 mode each byte of the pattern colour table provides a foreground and background colour for the corresponding byte in the pattern generator table. Byte 0 of the pattern colour table maps directly onto byte 0 of the pattern generator table. Byte 1 onto byte 1. Byte 2 onto byte 2 and so on. Assuming the pattern generator table is located at 0K and the pattern colour table at 8K, to discover which pattern generator byte is coloured in by which colour byte all we need to do is add an offset which is equal to the distance between the two tables to find the desired VRAM address. In this case the offset is 8192 Dec (£2000 Hex).

The mapping scheme below shows how the coloured patterns are actually mapped onto the display screen.

^ (Illus 3 – Diagram of mapping scheme)

In order to access this mapping scheme, the pattern name table needs to be initialised correctly. In order for the first pattern character block to map onto the first character cell on the screen and the second pattern character block to map onto the second character cell on the screen and so on, screen position 0 must contain 0, screen position 1 must contain 1, up to screen position 255 which must contain 255.

This however only accounts for the top third of the screen, what of the other two thirds?

Internally the VDP segments the pattern generator, pattern colour and pattern name tables, into three equal blocks of 2048, 2048 and 256 bytes respectively. Character labels in the upper third of the screen automatically correspond to the character patterns in the upper third of the pattern generator table. Labels in the middle third of the screen correspond to the character patterns in the middle third of the pattern generator table, and labels in the lower third of the screen correspond to the character patterns in the lower third of the pattern generator table.

The second and third blocks of the screen are mapped onto the related character patterns in the pattern generator table on the basis of position, therefore, in order to fill the second and third parts of the screen with the correct labels we load position 256 with 0 through to 511 with 255, and position 512 with 0 through to 255.

The correct start label values for the second and third block of the screen of 256 and 512 respectively are added by the VDP.

Assuming that we have located our pattern name table at 15K, then the assembler code routine which will set up the screen correctly is shown below. The routine uses the VRAM I/O routines described in section 1.5 and 1.6

```

;
; INISCR      INITIALISE PATTERN NAME TABLE (DISPLAY SCREEN)
;
;            NO PARAMETERS REQUIRED ON ENTRY
;            NO REGISTERS AFFECTED ON EXIT
;
INISCR:      PUSH      AF          ;Save Acc and flags
             PUSH      BC          ;Save BC register pair
             PUSH      DE          ;Save DE register pair
             PUSH      HL          ;Save HL register pair
             LD        HL,768      ;Set up loop counter to equal
                                   ;768 – This is the size of the
                                   ;display screen
             LD        C,0         ;Set byte to be output to each
                                   ;sequential screen position to zero
             LD        DE,15360    ;Load DE with start address of screen
             CALL     VSETOT       ;Set VDP write to VRAM pointer
                                   ;to point to start of screen
;
INISC1:      CALL     VDOUTP       ;Output display byte to screen
             INC      C            ;Increment display byte
             DEC      HL           ;Decrement loop counter
;
             LD      A,H          ;If loop counter < > 0 then
             OR      L            ;goto INISC1 else drop through
             JR      N,INISC1     ;to INISC2
;
INISC2:      POP      HL          ;Get old HL register pair
             POP      DE          ;Get old DE register pair
             POP      BC          ;Get old BC register pair
             POP      AF          ;Get old Acc and flags
             RET                ;Return to calling routine

```

Text Mode 2.5

Text mode is either decimal or hex, alphanumeric, or the low byte of a label.

eg: - <label>: DB <dec/hex No.> and/or “<string>” and/or <label>

NB: Using a label will generate an out of range error but ignoring this error leaves the low byte of the label 8 (Vertical) pixel size.

The character patterns can be dynamically changed to give any number of character patterns limited by the amount of storage space in Z80 RAM for the extra pattern libraries.

The display colours can be any one of 15 including transparent for the backdrop colour, and any one of 15 including transparent for the text colour.

The VDP is in text mode, when mode bits M1 = 1, M2 = 0 and M3 = 0 (see section 2.2 register 1 bits 3 and 4).

When the VDP is first initialised into text mode the VRAM is organised as shown in the table below.

VRAM sub-block	Length in bytes
Pattern generator table	2048
Pattern name table	960

Text mode VRAM arrangement in BASIC has already been discussed in section 2.4, because in BASIC it has been designed to be an integral part of the tabling for graphic mode 2.

It can be seen however, that the text pattern generator table as BASIC sets it up is only 1K long. This will only allow you to have a maximum of 128 different patterns in this mode. As explained in section 2.4 concerning the reduction in size of the sprite generator table, this is a compromise measure in order to have both graphics two mode and text mode resident in VRAM at the same time.

Sprites are not available in this mode because it is concerned only with the use of text type characters.

Text mode has many advantages in that it is very compact. A text mode set-up only occupies 3K of VRAM. It is possible to build up several different text libraries in VRAM and several screens, and switch to a complete new text set-up by changing only two VDP write-only registers.

Alternatively it is possible to have one text library held in the text pattern generator table, and to have up to 14 completely separate display screens, which can be changed to a new screen by changing only one VDP write-only register (Register 2), in a time interval of about 11 micro-seconds. This would be very useful for building up animated displays.

The text pattern generator table is 2048 bytes long and is split into 256 text patterns, each of which is 8 bytes long. Since each text position on the screen is six pixels across, the least significant bits (ie bits 1 and 0) of each text pattern byte are ignored. Each block of eight bytes in the text pattern library define a text pattern in which the 1's take on the text foreground colour, while the 0's take on the background (or backdrop) colour. These colours are chosen by loading VDP register 7 as described in section 2.2 example 1.

Any one of any of the patterns held in the text pattern library can be displayed in any position on the current text display screen, simply by loading the value of the desired pattern number into the appropriate position in the text display screen area.

Assuming we had a standard ascii character set held in our pattern library and that the text display screen was located at 7K, then we could use the simple routine below to print out strings to this screen. The routine uses the VRAM I/O routines as described in section 1.

The ascii values of the characters we wish to output will provide the correct pattern number to be loaded into the text display screen area to extract the correct pattern shape.

```

;
;STRING    -PRINT A STRING ROUTINE
;
;          VDP SHOULD BE IN TEXT MODE ON ETYR
;          START ADDRESS OF STRING POINTED TO BY HL
;          SCREEN ADDRESS TO WHICH STRING MUST GO IN DE
;          STRING MUST BE DELIMITED BY A '$'
;
;          NO OTHER PARAMETERS REQUIRED
;          NO REGISTERS AFFECTED EXCEPT HL WHICH POINTS
;          TO '$' CHARACTER ON EXIT
;
MESSAGE:  DB    'This is an ascii string$'
;
START:    LD    HL,MESSAGE    ;Set HL to point to start address of
;                          ;text string
          LD    DE,7144      ;Set DE to point to start address at
;                          ;which the text string is to be output
          CALL STRING        ;Output string to screen
          RET                ;Return to calling routine

;
STRING:   PUSH AF            ; Save Acc and flags
          PUSH BC            ; Save Bc register pair
          CALL VSETOT        ; Set VDP write to VRAM pointer to start
;                          ;address for text output

STRIN1:   LD    A,(HL)       ;Load text byte into Acc
          CP    '$'          ;Test to see whether if character is
;                          ;a '$'
          RET  Z              ;If true then return to calling routine
;                          ;else drop through to line STRIN2

STRIN2:   LD    C,A          ;Load text byte into output register
          CALL VDOUTP        ;Output text byte to VRAM
          INC  HL             ;Increment text pointer to point to next
;                          ;byte in text string
          JR   STRIN1        ;Go back to STRIN1 and do it again

```

Graphics 1 Mode 2.6

Graphics 1 mode is the other graphic mode available on the MTX and is not normally allowed from BASIC. It can be enabled with ease in assembler, and does have certain useful features. The features it provides are summarised in the table below.

Screen	768 character positions (24 rows by 32 columns) with up to 256 unique characters at any time in an 8 (Horizontal) by 8 (Vertical) pixel size.
--------	---

The graphics pattern display colours can be any two of 16 (including transparent) sectioned into groups of 8 characters.

The backdrop or border colour can be any one of 16 (including transparent).

The character patterns can be dynamically changed to give any number of character patterns limited by the amount of storage space in Z80 RAM for the extra pattern libraries.

Sprites are available up to a maximum of 32.

The VDP is in Graphics 1 mode, when mode bits M1 = 0, M2 = 0 and M3 = 0 (see section 2.2 register 1 bits 3 and 4).

One of the major advantages of this mode is that it is very compact, requiring a maximum 2848 VRAM bytes for a complete set-up. Yet it is possible with some care to create graphics effects very similar to those available in graphics 2 mode.

Like text mode, it is also possible to build up several different pattern libraries in VRAM and several screens, and switch to a complete new graphics 1 mode set-up by changing only two VDP write-only registers.

And again like text mode it is possible to have one pattern library held in the pattern generator table, and to have up to 14 completely separate display screens, which can be changed to a new screen by changing only one VDP write-only reg A,E ;Get data byte

```
OUT (2),A           ;Output data byte
LD A,D             ;Get control byte
AND 7              ;Set up correct control bits
OR 128             ;Set bits 6,0w
```

VRAM sub-block	Length in bytes
Pattern generator table	2048
Pattern colour table	32
Pattern name table	768

The pattern generator table is 2048 bytes long and is split into 256 graphics patterns, each of which is 8 bytes long.

Each block of eight bytes in the pattern generator library define a graphics pattern in which the 1's take on the foreground colour assigned to its, while the 0's take on the appropriate background (or backdrop) colour.

These colours are chosen by loading the correct byte of the pattern colour table with the appropriate colour byte in a very similar way to that used to load VDP register 7 as described in section 2.2 example 1.

The difference is that these colour control bytes are held in VRAM and do not need a register type write to set them up, but would be manipulated using the VRAM I/O routines described in section 1.

The mapping arrangement for pattern generator character shapes to the appropriate colour bytes is shown in the table below.

Pattern numbers	Pattern colour bytes	Pattern numbers	Pattern colour bytes
0 to 7	0	128 to 135	16
8 to 15	1	136 to 143	17
16 to 23	2	144 to 151	18
24 to 31	3	152 to 159	19
32 to 39	4	160 to 167	20
40 to 47	5	168 to 175	21
48 to 55	6	176 to 183	22
56 to 63	7	184 to 191	23
64 to 71	8	192 to 199	24
72 to 79	9	200 to 207	25
80 to 87	10	208 to 215	26
88 to 95	11	216 to 223	27
96 to 103	12	224 to 231	28
104 to 111	13	232 to 239	29
112 to 119	14	240 to 247	30
120 to 127	15	248 to 255	31

It can be seen that for each block of eight patterns in the pattern generator library the colour is determined by one byte of the pattern colour cable. This means that to obtain a wide range of differently coloured graphic patterns in this mode may require some thought.

Any one of any of the patterns held in the pattern generator library can be displayed in any position on the current graphics display screen, simply by loading the value of the desired pattern number into the appropriate position in the graphics display screen area.

Assuming that we had set up the pattern colour table bytes as below:

Pattern colour table byte 0	=	£F1 (Hex), 241 (Decimal)
		White on black
byte 1	=	£41 (Hex), 65 (Decimal)
		Blue on black

The graphics display screen was located at VRAM address £C00 (Hex), 3096 (Decimal) or 3K.

The pattern generator library was located at VRAM address £0 (Hex), 0 (Decimal), or 0K.

Pattern generator table block	0	=	Blank pattern
	8	=	Diamond shape

Then a routine which will blank the screen out and draw a border around the outside of the screen is listed below. This routine makes full use of the VDP I/O routines listed in section 1.


```

;
; address held in DE HL number of
; times:
; Return to calling routine
RET

;
; BOTTOM
LD DE,£EA1
CALL VSETOT
; Set the VDP write to VRAM
; pointer to the start of the bottom
; line of the graphics display screen
LD C,8
; Load the border character into
; the output register
LD HL,32
CALL OUTBLK
; Loop counter = 32 (Decimal)
; This routine will output the byte
; held in C to VRAM from the start
; address held in DE HL number of
; times
RET
; Return to calling routine

;
; SIDES:
LD DE,£C20
; Load DE with the start address of
; the second line of the display screen
; Loop counter = 22 (Decimal)
; Save old BC register pair
; Set the VDP write to VRAM pointer
; to the start address contained within
; DE
; SIDES1:
LD C,8
CALL VDOUTP
; Load C with the border character
; Output border character to display
; screen (Left hand side)
LD HL,31
ADD HL,DE
EX DE,HL
CALL VSETOT
; Increment DE by 31 ie move the
; VDP write to VRAM pointer
; to the right hand side of the screen
; Rest the VDP write to VRAM pointer
; to the start address contained within
; DE
CALL VDOUTP
; Output the border character still
; retained in C to the display screen
; (Right hand side)
INC DE
; Increment the VDP write to VRAM
; reference pointer by one to move
; along character position and down
; one line back to the left hand side
; of the screen
POP BC
DJNZ SIDES1
; Get old BC register pair
; Decrement loop counter and if loop
; counter < > 0 then goto SIDES1
; else drop through to SIDES2.
; Return to calling routine
; SIDES2:
RET

;
; OUTBLK:
CALL VDOUTP
; Output data byte held in C to
; VRAM to address pointed to by the
; VDP write only pointer
DEC HL
LD A,H
; Decrement loop counter held in HL
; if loop counter < > 0 then goto

```

```

                OR          L          ;OUTBLK else drop through to
                JR          NZ,OUTBLK ;to OUTBL1
OUTBL1:        RET          ;Return to calling routine

```

Multicolour Mode 2.7

Multicolour mode is not normally allowed from BASIC, but like graphics 1 mode it can be enabled with ease in assembler. Although I cannot see much use for it, it is available on the machine and so I will attempt to describe it.

The features it provides are summarised in the table below:

Screen	An unrestricted 48 row by 64 column display consisting of blocks of 4 (Horizontal) by 4 (Vertical) pixels in any one of 15 colours plus transparent.
	The character pattern colours can be dynamically changed to provide colour animated displays.
	The backdrop or border colour can be any one of 15 plus transparent.
	Sprites are available up to a maximum of 32.

The VDP is in multicolour mode, when mode bits $M1 = 0$, $M2 = 1$ and $M3 = 0$ (see section 2.2 register 1 bits 3 and 4).

Multicolour mode occupies a total of 1728 VRAM bytes in a complete set-up divided into two tables, but because the tables are not contiguous and begin on even 1K and 2K boundaries a total of 3K is needed. 768 bytes are used for the name table and 1536 bytes (24 rows by 32 columns by 8 bytes per pattern position), for the pattern generator table.

Like the graphic modes, multicolour mode has a screen consisting of 768 character positions. The character label value contained within any one of the positions on the display screen does not point to a character shape because this is always a 2 (Horizontal) by 2 (Vertical) block of 4 (Horizontal) by 4 (Vertical) pixels all set to 1,s. Instead it points to a colour reference held within what would normally be the pattern generator table. As in other modes each cell within the pattern generator table consists of eight bytes. Multicolour mode only uses two of these bytes within each pattern generator cell for each label on screen.

These two bytes specify four colours, each colour being related to each of the 4 by 4 pixels blocks within a character position on screen. The four MSB,s of the first byte define the colour of the upper left quarter of the multicolour pattern. The four LSB,s of the first byte define the colour of the upper right quarter. The second byte similarly defines the lower left and right hand corner of the multicolour pattern.

This is shown in the diagram below:



^(Illus 4 – ABCD diagrammed from above)

The location of the two bytes within the eight byte segment of the pattern generator table pointed to by the character label value held in any one of the display screen character positions is dependent on the screen position in which the character label is held. This is also elaborated to some extent by the diagram below.

For names in the top row of the display screen (ie values 0 to 31), the first two colour bytes of the pattern generator cell are accessed. For the second row of the display screen (ie values 32 to 63), the second two colour bytes of the pattern generator cell are accessed. The next row of the screen uses the fifth and sixth bytes and the next row uses the seventh and eighth . This series repeats for the remainder of the screen.

Pattern Generator Cells	Screen Rows	Pattern Generator Bytes
0 to 31	0	0 and 1
32 to 63	1	2 and 3
64 to 95	2	4 and 5
96 to 127	3	6 and 7
128 to 159	4	0 and 1
160 to 193	5	2 and 3
(X) to (X+31)	(X DIV 32)	(X DIV 16 MOD 8 and (X DIV 16 + 1) MOD 8

NB: X = start location of each of the display screen rows relative to the top left hand of the screen.

Points to look out for 2.8

One of the major points to look out for when manipulating the VDP chip through any of the four different modes is a direct result of the action of BASIC.

BASIC services VRAM at all sorts of odd times using various interrupts. If you are using the same modes as BASIC uses there will be no problems if you are using them in the same way as BASIC does. If however, you are using and creating your own VDP modes or are using the same modes with different VRAM table set-ups you will have to choose between having BASIC around or not, else BASIC and your code will conflict and you will end up with garbage.

In most cases if you are doing anything that BASIC would not normally do you will have to use assembler. BASIC can easily be disabled by using a DI instruction at the start of your code, and an RETI instruction at the end. If you have modified the contents of the VDP registers drastically it would be far better to perform a JP £0000 as your last instruction as a warm boot back to BASIC.

No matter what processes you perform in Z80 ram, as long as you do not corrupt the VDP registers, or perform a BASIC reboot, or perform a system reset the integrity of VRAM will always be maintained.

You can switch from one VDP mode to another in mid processing operation by performing a VDP write only register change. Some very interesting effects can be obtained by having graphics 1 mode tables and text mode tables on compatible VRAM start boundaries and then switching from one to the other.

All direct VDP operations are associated with ports 1 and 2 only.

VDP and CTC Interrupts – Introduction

3.1

At the end of each active display scan which is about every 1/50 of a second, the VDP chip will stop all screen processing and perform other tasks. It is possible to set up VDP interrupt servicing routines which will start at the end of each active display scan, and finish before the next one has begun.

This routine is a period of time when data bytes written to the screen or blanked on-screen can be processed in a glitch free manner.

Bit 5 of register 2 must be set to a 1 to allow VDP interrupts to take place. This should be done during VDP write only register set up at the start of your code.

In addition to this the use of standard IM 2 interrupts has been enhanced using the CTC chip on board the MTX, to allow the user to set up vector tables on an 8 byte boundary within a page instead of a page boundary alone leading to a much greater flexibility and ease of use in vector table positioning.

This can be a complex subject for some computer users who are new to the subject, but it is well worth while persisting with this section. The results of mastering interrupts make the effort involved small in comparison.

The advantages of using VDP interrupts are:

1. A drastic improvement in the quality of animated displays. Reducing or removing glitches created when transferring bytes to or from the display screen area.
2. It gives a second optional clock and can be used for timing and compact delay loops.

Programming the CTC

The Zilog CTC counter Timer Circuit handles all interrupts on the MTX including the Video Display Processor (VDP) interrupt. The following is intended only as brief outline concerning CTC operation. For more details refer to Zilog's Z8430 CTC Counter/Timer Circuit product specifications.

The CTC is capable of generating mode 2 interrupts from any of its 4 independently programmable channels. It is capable of acting as either a timer or counter, working on an external clock. The port numbers, CTC channel numbers and functions are as follows:

Port Number	Channel	Function
08	ch0	VDP interrupt line
09	ch1	4 MHz system clock /13
0A	ch2	4 MHz system clock /13
0B	ch3	cassette edge input

The first word to write out to the channel being programmed is the channel control word this is made up as follows:

Bit	Value	Function
B7	1	Interrupt enabled
	0	Disable interrupt
B6	1	counter mode
	0	timer mode
B5	1	prescaler of 256 *
	0	prescaler of 16
B4	1	trigger on rising edge
	0	trigger on falling edge
B3	1	clk/trig pulse starts timer *
	0	start on receiving time constant
B1	1	software reset
	0	continued operation
B0	1	control word
	0	vector word

Using Timer Mode Only

If bit 2 is set then the channel will consider the next byte it receives to be a time constant. Setting bit 1 causes the channel to stop what it is doing and accept the next set of parameters it will also require a new time constant.

The interrupt vector word is identified by a zero in bit 0. The 5 most significant bits form the 5 most significant bits of the interrupt vector provided by the chip on interrupts (mode 2). Bits 1 and 2 are set according to the channel generating the interrupt and bit 0 is always zero.

B2	B1	Channel
0	0	0
0	1	1
1	0	2
1	1	3

The interrupt vector table must lie on an 8 byte boundary. This table normally sits at FFF0 hex. A working example set up of an IM 2 vector table is detailed in section 3.2

Note that the ctc channels were reset twice, this is because one or all of the channels may be expecting a timer constant and hence misunderstand the first reset program word sent to the channel.

In the service routine the VDP status register is read (IN (2),A). This is to acknowledge and reset the VDP interrupt as well as re-initialise the read write cycle of the VDP.

How to Generate VDP Interrupts 3.2

Setting up a VDP interrupt servicing routine is essentially a three stage process which is detailed in parts 1 to 3. Part 4 is a listing of the code which will perform the complete task.

Before this section of code is executed it is vital to inform the VDP chip that it is going to be expected to perform user defined interrupts. This is done by ensuring that when the VDP write only registers are set up, bit 5 is register 1 is at 1. This is the VDP interrupt enable bit.

Part 1

VDP interrupts are set up and accessed through the normal Z80 IM 2. This mode is tied into the CTC chip and enables four channels to be made available to the user on ports 8, 9, 10 and 11 referred to as CTC channels 0, 1, 2 and 3 respectively.

The first task that needs to be done is to shut off any existing interrupts on board the ctc chip which you do not require. The code below will switch off all ctc interrupts on channels 0, 1, 2 and 3.

```

KILLcTc:   LD    B,2           ;Loop counter = 2
           LD    A,3           ;Ready reset CTC channel byte in
                                           ;Acc ready for output to ports 8, 9, 10 & 11
KILLloop:  OUT   (CTC),A       ;Ouput reset byte to channel 1 port 8
           OUT   (CTC+1),A     ;Ouput reset byte to channel 2 port 9
           OUT   (CTC+2),A     ;Ouput reset byte to channel 3 port 10
           OUT   (CTC+3),A     ;Ouput reset byte to channel 4 port 11

```

```

DJNZ KILLloop          ;Decrement loop counter and if loop
                        ;counter < > 0 then goto KILLloop else
                        ;drop through to next section of code.

```

It can be seen in the above code that the reset byte 3 is being written to the ctc chip twice. This is because the CTC may expect the next byte input to it to be a time constant therefore a re-write will eliminate these.

Part 2

After executing the code above, the next section of code (listed below), selects interrupt mode 2, sets up the interrupt vector table by loading the high byte of the vector table start address into the I register and the low byte of the vector table start address into channel 0 of the ctc chip, then loading the interrupt servicing routine start address into vector table bytes 0 and 1.

The most important point to note at this stage is that the two byte vector table start address selected by the value in I (High byte), and the value output to CTC channel 0 (Low byte), MUST point to a fixed 0 byte boundary address.

The assumption being made is that the start address of the vector table is £8180 (Hex).

```

SETupCTC: DI           ;Disable all existing interrupts
            IM    2     ;Select interrupt mode 2 (IM 2)
            LD    A,£81 ;Load the high byte of the vector
                        ;table start address into the
            LD    I,A    ;page select register
            LD    A,£80 ;Select low byte vector table
                        ;start address on 8 byte boundary
            OUT   (CTC),A ;within the page pointed to by the
                        ;contents of I.
            LD    HL,VDPout ;Select the start address of the actual
                        ;servicing routine. In this case it is
                        ;indicted by the label VDPout
            LD    (IJtable),HL ;Load the interrupt servicing routine
                        ;start address into vector table bytes
                        ;0 and 1.

```

Part 3

The final stage is to re-enable the ctc chip interrupt on channel one, and clear the vdp interrupt flag on board the vdp chip, by performing a read of the vdp read only register (signified by variable VDPRGO which = 2).

Each time the read only register is read, the interrupt line is reset on board the vdp chip.

NB:

1. The read only register on board the vdp chip must be read each time you exit from your interrupt routine
2. The maximum duration of your interrupt routine MUST NOT be longer than 1/50 second (20,000 microseconds).

```

SETupINT: LD    A,0C5H           ;Send CTC bytes to CTC chip to
          OUT   (CTC),A         ;'wake' it up and prepare CTC
          LD    A,1             ;chip prior to beginning execution of
          OUT   (CTC),A         ;VDP interrupts
          IN    A,(VDPRGO)      ;Clear VDP interrupt flag held in VDP
          EI                    ;read only register on port 2
          RETI                   ;Enable new interrupt system
          RETI                   ;Return to calling routine and exit from
                                ;interrupt set up routine.

```

Part 4

Complete source listing

```

;
;VARIABLES SECTION
;
CTC      EQU 8
VDPRGO   EQU 2
;
;TABLES SECTION
;
IJTABLE:      0,0,0,0,0,0,0,0   ;THIS TABLE MUST BE LOCATED
                                ;ON AN 8 BYTE BOUNDARY
                                ;In this case address £8180
;
;VDP INTERRUPTS
;
KILLcTc:  LD    B,2             ;Shuts off all CTC channels
          LD    A,3
KILLloop: OUT   (CTC),A
          OUT   (CTC+1),A
          OUT   (CTC+2),A
          OUT   (CTC+3),A
          DJNZ KILLloop
;
SETupCTC: DI
          IM    2
          LD    A,$81
          LD    I,A
          LD    A,£80
          OUT   (CTC),A
          LD    HL,VDPout
          LD    (IJtable),HL
;
SETupINT: LD    A,0C5H
          OUT   (CTC),A
          LD    A,1
          OUT   (CTC),A
          IN    A,(VDPRGO)
          EI

```

```

                RETI
;
VDPout:        DI                ;Disable all interrupts
;
                (start of your routine)
(you MUST save all registers which will be affected by the servicing routine at this
point).

```

```

                IN    A,(VDPRGO)    ;Has the VDP reached the end of
                BIT   7,A           ;the current active display
                JR    NZ,VDPou1     ;scan – Indicated by bit 7 of the VDP
                                        ;read only register = 1
                                        ;If condition is true then goto VDPou1
                                        ;and begin VDP servicing routine else

```

(Retrieve all saved registers)

```

                RET1                ;exit interrupt routine and return to
                                        ;calling routine
;

```

```

VDPou1:        (VDP servicing begins)
                “

```

```

                (Insert whichever section 3.3)
                (routines you are using at this point)

```

```

                IN    A,(VDPRGO)
                (end of your routine)
                “

```

(Retrieve all saved registers)

```

                RET1

```

Using VDP Interrupts

3.3

The most important point to note when using VDP interrupts is that whatever code you write to go in them MUST NEVER take longer than 1/50th of a second (20,000 microseconds).

All of the following sections of code are designed to be placed in the VDP servicing routine described in section 3.2 at the appropriate point.

As mentioned in section 3.1, the routine detailed in section 3.2 can be used to generate reasonably accurate and useful clocks for time dependent processes or delays within say games programs. Below is listed a section of code which will generate a hundred hour clock and return the results in a 6 byte table called CLOCK.

The facilities it offers are:

1. The contents of CLOCK can be examined at any time to return a hundred hour value for
 - A) Hours - Bytes 0 and 1
 - B) Minutes - Bytes 2 and 3
 - C) Seconds - Bytes 4 and 5

An important point to note is that the values of the digits returned by this routine are ASCII values.

2. Every 1/50th of a second the variable ONE50 is set to one and can be used for delay loop timing.
3. By using a variable called CLRCLK you have the following:

- A) CLRCLK = 0 – No action taken
- B) CLRCLK = 1 – Clear CLOCK. Time is '00 00 00'
- C) CLRCLK = 2 – CLOCK is set to time contained within 6 byte table TIMSET.

```

;
;SEVCLK – INTERRUPT SERVICING ROUTINE SECTION
;   IM 2 – POINTER TABLE IJTABLE SET UP USING CODE
;   LISTED IN SECTION 3.3
;
;
;   ROUTINE MAINTAINS 6 DIGIT HUNDRED HOUR CLOCK
;
;   USES REFERENCE TABLE HELD IN CDATA (CLOCK DATA)
;
;   ALSO CONTROL VARIABLE CLRCLK
;
;   NO PARAMETERS REQUIRED ON ETRY
;   REGISTERS AFFECTED ON EXIT ARE AF, BC, DE AND HL
;
;
CLOCK:
HOURS:   DB      30H,30H      ;(see description above)
MINS:    DB      30H,30H      ;(see description above)
SECS:    DB      30H,30H      ;(see description above)
WIDGET:  DB      30H          ;1/50th of a second widget
;counter
ONE50:   DB      0            ;see description above
CDATA:   DB      '99','59', '59' ;set up clock counter reference
;data
TIMSET:  DB      '1', '1', '3', '4' ;6 byte table used to reset
;table CLOCK to some specific
;time when CLRCLK = 2

CLRCLK:  DB      0            ;see description above
;
SEVCLK:  LD      A,1          ;set 1/50th second counter to 1
;LD      (ONE50),A
;LD      A,(CLRCLK)         ;Check value of CLRCLK.
;CP      0                  ;and drop through to SEVCL2
;JP      NZ,SEVCL5         ;else goto SEVCL5
SEVCL2:  LD      DE,CLOCK+6   ;Increment clock by 1/50th
;LD      HL,CDATA+6        ;of a second using one hundred
;LD      B,7               ;hour reference values in table
SEVCL3:  LD      A,(DE)       ;CDATA
;CP      (HL)
;JR      C,SEVCL4
;
SEVCL3A: LD      A, '0'-1
SEVCL4:  INC      A

```

```

LD      (DE),A
JP      C,SEVCL7
DEC     DE
DEC     HL
DJNZ    SEVCL3
JP      SEVCL7      ;Goto routine exit at this point
;
SEVCL5: CP      2
JP      NZ,SEVC5D  ;Has set a time option been
                    ;selected at this point
                    ;If condition true then drop
                    ;through SEVC5A else goto
                    ;SEVC5D
;
SEVC5A: LD      HL,TIMSET      ;Set clock to specific time held
                    ;in table TIMSET and then
                    ;goto SEVCL7 and exit routine
LD      DE,CLOCK
LD      BC,7
LDIR
LD      A,0      ;Reset option select control
LD      (CLRCLK),A ;variable to zero
JP      SEVCL7
;
SEVC5D: LD      HL,CLOCK      ;Clear clock option selected at
LD      A,30H      ;this point here
LD      B,7      ;Fill table CLOCK with ASCII 0,s
SEVCL6: LD      (HL),A      ;and then exit routine
DJNZ    SEVCL6
LD      A,0      ;Reset option select control
LD      (CLRCLK),A ;variable to zero
;
SEVCL7: (End of routine)

```

The following section of code illustrates the use of the above routine. It uses a table called STATIM (Start time), which contains the start time for this section of code.

```

;
SAMPLE: LD      A,1      ;Select clear clock option
LD      HL,STATIM      ;Transfer clock start time
LD      DE,TIMSET      ;from table STATIM to table
LD      BC,6      ;TIMSET
LDIR
LD      A,2      ;Select clock reset option
LD      (CLRCLK),A      ;and update clock with new value
                    ;contained within table TIMSET
RET      ;Return to calling routine
;

```

Another useful routine which can be used within the code described in 3.2 is a random number routine.

Each time you require a random number (8 bit value), it is only necessary to look at the location RND to extract a new value. This will be constantly updated under interrupt by the interrupt routine RANDOM.

A point to note is that the random value contained within RND will only be auto-updated every 1/50th of a second. If this is too slow for some applications then it will do no harm to call the routine random independently of the interrupt routine using a standard Z80 'CALL RANDOM' instruction.

The assumption made is that prior to setting under the VDP interrupt routine a routine called SETRND is called. This is listed below. It sets up a random seed within an 8 byte table RNDMEM for use by the routine RANDOM.

```

;
; SETRND-SET UP SEED WITHIN 8 BYTE TABLE RNDMEM FOR USE
;           BY ROUTINE RANDOM
;           NO REGISTER SET UP REQUIRED ON ENTRY
;           NONE AFFECTED ON EXIT
;
RNDMEM:  DB      0,0,0,0,0,0,0,0
;
SETRND:  PUSH    AF           ;Save Acc and flags
         PUSH    BC           ;Save BC register pair
         PUSH    IY          ;Save IY index register
         LD      B,5          ;Write five seed values into table
         LD      IY,RNDMEM    ;RNDMEM extracted from the Z80
SETRN1:  LD      A,R          ;refresh register
         LD      (IY+0),A     ;Extra 3 bytes in table RNDMEM
         INC     IY           ;used by routine RANDOM
         DJNZ   SETRN1
         PDP    IY           ;Retrieve old IY index register
         PDP    BC           ;Retrieve old BC register pair
         PDP    AF           ;Retrieve old Acc and flags
         RET                ;Return to calling routine

```

The code for the random number routine follows

```

;
; RANDOM-RETURN 8 BIT RANDOM VALUE IN VARIABLE RND
;           USES TABLE RNDMEM
;           NO REGISTER SET UP REQUIRED ON ENTRY
;           NONE AFFECTED ON EXIT
;
RND:     DB      0           ;See description above
;
RANDOM:   PUSH    AF           ;Save Acc and flags
         PUSH    BC           ;Save BC register pair
         PUSH    IY          ;Save IY register pair
         LD      B,8          ;Loop counter = 8
         LD      IY,RNDMEM    ;Set pointer to start of random
                                ;seed table
RAN0:    LD      A,(RNDMEM+2) ;Extract seed value
         SRL    A             ;Shuffle seed table
         SRL    A
         SRL    A
         XOR    (IY+4)
         RR     A

```

```

                RL      (IY+0)
                RL      (IY+1)
                RL      (IY+2)
                RL      (IY+3)
                RL      (IY+4)
                DJNZ     RAN0          ;Decrement loop counter and if
                                        ;loop counter < > 0 goto RAN0
LD      A,(IY+0)                    ;else drop through to next
                LD      (RND),A      ;statement and update value of RND
                PDP     IY           ;Retrieve saved IY register pair
                PDP     BC          ;Retrieve saved BC register pair
                PDP     AF          ;Retrieve saved Acc and flags
                (End of routine)

```

As mentioned in section 3.1, VDP interrupt routines can be used to create glitch free screen updates.

Listed below is a routine which will fill a text mode display screen with a selected byte (It could be a blanking byte or any other), simply by selecting the appropriate byte and loading it into a variable called FILTEX.

If FILTEX contains 255 no action is taken and therefore this is the only byte which cannot be sent to the screen using this routine.

The assumptions made when using the routine listed below are :

1. That the VDP is in text mode
2. The start address of the screen (pattern name table), is £1800 (Hex) or 6K
3. That we are using the VDP I/O routines described in section 1.1

```

;
;
;TEXTMOD-FILL A TEXT MODE SCREEN (960 BYTES LONG) WITH A
;      SELECTED BYTE HELD IN VARIABLE FILTEXR
;      IF FILTER = 255 THEN NO ACTION TAKEN
;
;
;      NO REGISTER SET P REQUIRED ON ENTRY
;      AF, BC, DE REGISTERS AFFECTD ON EXIT
;
;
;      THE TWO BYTE VARIABLE SATMSC
;      (START ADDRESS TEXT MODE SCREEN) MUST BE SET TO START
;      OF TEXT PRIOR TO ENTRY TO THIS ROUTINE
;
SATMSC:  DW      £1800          ;See description above
FILTEX:  DB      0             ;See description above
;
TEXTMOD:  LD      A,(FILTEX)    ;If variable FILTEX = 255
          CP      255          ;then no action taken – goto TEXMO4
          JR      Z,TEXMO4     ;and exit routine
          ;else drop through to TEXM01 and
          ;perform screen update
;
;
TEXTMO1:  LD      DE,(SATMSC)   ;Set write to VRAM pointer to start
          CALL    VSETOT       ;of text display screen
          LD      C,A          ;Load contents of variable FILTEX

```

```

;into the write to VRAM output
;register
;Loop counter = 960
TEXMO2: LD      DE,960
        CALL   VDOUTP ;Output fill screen byte to screen
        DEC    DE      ;Decrement loop counter and if
        LD     A,D     ;loop counter <> 0 then goto TEXMO2
        OR     E       ;else drop through to TEXMO3
        JR     NZ,TEXMO2

TEXMO3: LD      A,255 ;Reset fill screen byte to 255
        LD     (FILTEX),A
TEXMO4: (End of routine)

```

Listed below is a section of code which demonstrates the use of the above routine. It is assumed that the pattern generator library has been loaded with a series of ASCII patterns for text mode use.

```

;
SAMPLE: LD      A,'O' ;Fill text screen with the symbol
        LD     (FILTEX),A ;'O'
        LD     A,'X' ;Then fill text screen with the
        LD     (FILTEX),A ;symbol 'X'
        JR     SAMPLE ;Then goto sample and do it again
;The update will occur so smoothly
;that these two characters will
;merge to form a ????. (diagram).

```

It is a simple matter and very convenient from a programming point of view to be able to perform sprite collision detection routines under VDP interrupt.

The section of code listed below has been designed so that you have the option of enabling the test or disabling the test through a variable called IONOFF (Impact On/OFF) :

When IONOFF = 0 - No action taken
 IONOFF = 1 - Impact test routine takes place

More important however, is the fact that it is now not necessary to keep repeatedly calling the impact routine from within your main code but only to look at a single variable TRUFAL:

When TRUFAL = 0 – No impact has taken place
 TRUFAL = 1 – Then sprite impact has taken place

The sprite impact detection routine is that described in section 4.4. It is important to note that if you are going to use this routine and because of the amount of time each sprite impact test takes, you are constrained to a maximum number of ????. sprites which you must not exceed.

It is assumed that a sprite attribute table delimiter has already been set up prior to setting up the interrupt routine using the BLKSPR routine also described in section 4.

8. Now add the value obtained by subtracting the original value of VAZERO (which you noted down), from the original value of NBTOP, BASTOP and BASTBO.
9. Set VAZERO to its original value
10. Exit panel and return to BASIC
11. Save the appended code

At this point you will have an appended section of code saved on tape.

Screen Output using RST 10

3.5

The ROM calls for screen output are all in the form of restart 10 calls. Following each of these calls is data which tells the ROM routine what to do. At first this is a little confusing as data is stored in the path of the program, but is in fact remarkable easy to use.

Writing ASCII characters to the Screen

The following RST 10 routine passes the ascii representations of registers B and C to the screen.

```
START:    LD BC,"TM"
          RST 10
          DB 192
          RET
```

A less trivial example would be the following:

```
START:    LD, E,8
          LD HL, DATA
LOOP:    LD A,(HL)
          LD B,0
          LD C,A
          RST 10
          DB 192; write BC token
          DEC E
          JR NZ, LOOP
          RET
DATA:    DB "MEMOTECH"
```

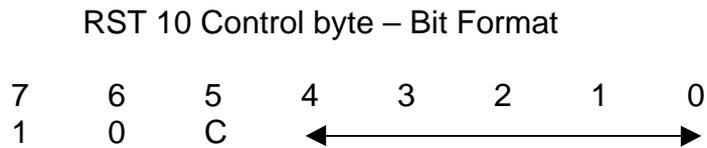
Try changing the line that loads B with zero to load B with the space character and see what happens.

Sending Messages to the Screen

To avoid the complication of the above routine we can send a complete string in the following way:

```
RST 10
DB £8D, "MEMOTECH LTD"
RET
```

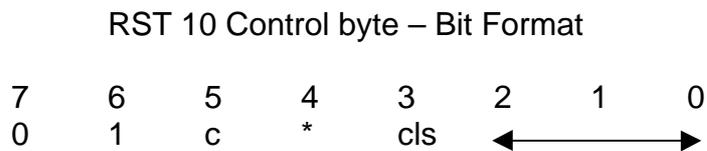
The byte following the RST 10 is made up in the following way:



Where bit 5 indicates that the routine should continue to interpret data after this instruction. n is the number of bytes in the string.

Virtual Screen and RST 10

The format for the virtual screen RST 10 instruction byte is:

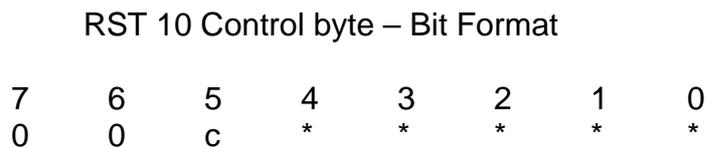


Where c is the continuation bit, n the virtual screen to be selected and cls gives the option to clear the screen.

* = don't care. (It can assume any value and has no operational effect when used in this mode).

One Byte Screen Write

This RST 10 call allows the transfer of single bytes to the screen.
Its format is:



Where c is the continuation bit

* = don't care. (it can assume any value and has no operational effect when used in this mode).

The following is an example using RST 10 and CALL £79, keyboard input.

```
START:   CALL £79
         JR Z,START
         LD C,A
         LD B,0
         RST 10
         DB 192
         CP 13
         JR NZ,START
         RET
```

The routine reads the keyboard and echoes the typed characters on the screen.

Control Codes and RST 10

In the ASCII character set there are 32 invisible characters before the first printable character (space). These invisible characters are called control characters. For example pressing both the control key and the 'G' key at the same time generates the bell code, character 7. These codes are extremely powerful in the MTX when used with RST 10.

Try the following example:

```
START:    LD B,160
LOOP:    RST 10
        DB £86, " * "
        DJNZ LOOP
LOOP:    RST 10
        DB £8B, 15,0, "***", £10,£10,£38
        DB £54,£10,£10,£10,£38
        CALL PAUSE
        RST 10
        DB £8B, 15,0,"**", 0,0,£10,£54
        DB £38,£10,£7C,£44
        CALL PAUSE
        JP LOOP1
        RET
PAUSE:   LD B,50
PAUSE1:  HALT
        DJNZ PAUSE1
        RET
```

The program works by first printing up a series of "***" characters and then redefining them to give an animation effect.

The following is a list of commands available through RST 10:

ASCII Code	Function
1	plot x,y
2	line x1, y1, x2, y2
3	cursor x,y
7	bell
10	line feed,cursor down
12	cls and home
13	vertical tab
13	carriage return
14	ctlspr p,x
15	genpat p,n,d0,d1,d2,d3,d4,d5,d6,d7
16	colour p,b
17	adjspir p,n,v
18	sprite n,p,yp,yp,xs,ys,col
19	movspir p,n,d
20	view dir,dis

21	insert key
22	delete key
23	back tab
25	tab key
27,65	attr p,state
27,89	crvs n,t,x,y,w,h,s
27,90	vs n
27,67	gr\$ x,y,b (result in work space)

Printer Output

All screen output can be redirected to either the RS232 or the centronics interface and hence the printer. To do this from basic type:

```
POKE 64143,DEV – Where DEV is    0    for screen
                                   1    for Centronics
                                   2    for RS232
```

Option 2 sets the list device. To divert screen output to the list device:

```
POKE 64885,1
```

This done simply in assembler, eg:

```
START:    LD A,1                    ;Select output to printer option
          LD (£FD75),A
          LD (£FA8F),A
```

Points to look out for

3.6

In almost every case that you use interrupts of the nature described in section 3.3 BASIC will be knocked out of action. This is the reason why there is no BASIC in this section because it cannot actually be used in conjunction with the listed code.

It is important to note that if you wish to exist from an interrupt servicing routine entirely for some time, then you must perform a DI instruction to eliminate any pending interrupts and then later an EI instruction to re-enable them.

If you wish to set up new interrupt servicing routines then you must perform a complete set-up as described in section 3.2

Sprites – Introduction

4.1

The best way to lead into this section is to define exactly what sprites are on the MTX. I am making the assumption that anybody who is reading this section will have seen sprites in action at some time, perhaps on an arcade game of some sort, or the MTX or some other computer.

Essentially a sprite on the MTX can be defined as below:

1. It is a special animation pattern which can be moved one pixel at a time, either vertically or horizontally dependent on or independently of the pattern background.
2. It can be coloured in any one of the 15 colours plus transparent available of the MTX independently of the pattern background.
3. It can display a pre-defined bit pattern shape held in VRAM by setting only one byte and can re-display other predefined shapes in a cyclic sequence to create true animated characters.
4. It can be 'bled in' off the top, left, bottom and right, of the display screen from 'behind' the border.
5. It can assume any one of several sizes and magnifications (see section 2.2 register 1).
6. Although the sprite X and Y coords, colours, shapes and a small amount of other information are created in software, all other aspects of sprite maintenance are performed under hardware.
7. Any sprite can be defined to pass in front of or behind any other sprite, and therefore, it is possible to create multilevel pattern overlaying.

The sections of VRAM which deal with the control and the definition of sprites are the sprite attribute table, and the sprite generator table (see section 4.2). Further additional control for size and magnification of sprites is determined by VDP write-only register 1 and is discussed in depth in section 2.2 Register 1. Additional useful information concerning sprites is contained in sections 4 and appendix B.

Sprite Attribute Table

4.2

The sprite attribute table is concerned with the control of sprites. As there are 32 sprite available on the MTX there are 32 control blocks within the sprites attribute table, each control block consisting of four bytes. This means that the sprite attribute table is 128 bytes (4 times 32 bytes), long.

The start address of the sprite attribute table is determined by the contents of VDP write only register 5, which locates the sprite attribute table on a 128 byte boundary (see section 2.2 register 5).

Each sprite on the MTX is assigned a priority by the hardware of the machine. This means that out of any two sprites the higher priority sprite will appear to pass in front of the lower, and the control block of which will occur nearer to the start of the sprite attribute table.

Sprite 0 being assigned to sprite attribute bytes 0,1,2 and 3, has the highest priority. Sprite 31 being assigned to sprite attribute bytes 124, 125, 126 and 127 has the lowest priority.

The diagram below shows exactly how the four control bytes within each control block of the sprite attribute table are arranged.

^ (Illus 5 – Sprite attribute table control block)

The first two bytes control the Y and X positions of the sprite onscreen.

The first byte indicates the vertical distance from the top of the screen in pixels, and is defined such that a value of -1 (ie £FF Hex), places the sprite at the top of the screen touching the backdrop area. A value of 0 will place the sprite one pixel below the top of the screen. A value of 1 will place the sprite one pixel below that and so on.

The second byte indicates the horizontal distance in pixels, from the left hand side of the screen, with a value of zero placing the sprite tight against the left hand border.

An important point is that all positioning of sprites on screen is determined from the upper left corner of any sprite, therefore if we use a 16 by 16 sprite as an example positioning the sprite at location 10,10 (Y/X):

^ (Illus 6 – 16 * 16 sprite pattern).

The third byte contains an 8 bit number (or name), which points to the sprite pattern held in the pattern generator table. The fourth byte of the sprite attribute table is concerned mainly with colour.

When the sprite pattern is defined in the pattern generator table, the 1,s of the pattern assume the colour held in the lower 4 bits of this byte. The 0,s within the sprite pattern are automatically set to transparent.

Byte four also performs one other function which is sometimes useful. Bit C is the Early Clock Bit. when it is set to zero it does nothing, but when it is set to one will shift the horizontal position of a sprite to the left by 32 pixels and allow a sprite to bleed in off the left hand side of the backdrop rather than the normal default right.

All of the control bytes within each control block of the sprite attribute table can be altered dynamically during the running of your program code. This means that within each control block you can:

1. Animate any of the sprites by changing the third byte
2. Create colour scintillating sprites by changing the fourth byte
3. Create moving sprites by updating the Y/X coords in control bytes one and two

As the information of sprite control is held in VRAM on pre-defined boundaries, you can create sprite impact routines with ease because for any sprite you can always refer to a fixed address independently of your program. It is also very easy to create sprite to background mapping routines because there is a straight forward and simple direct relationship between them.

The code and techniques which will perform the above mentioned two tasks are described in depth in section 4.4.

Setting up the Sprite Attribute Table

The start address in VRAM of the sprite attribute table is determined by the contents of VDP write only register 5, and is a fixed address. Given this fact and that there is a proportional relationship between any sprite number and its position in the sprite attribute table, it is easy to handle and manipulate sprites.

An important point to realise is that the contents of the sprite attribute table like any other part of VRAM could assume any indeterminate value when the machine is switched on and therefore must be cleared prior to use.

Clearing the sprite attribute table is not just a simple matter of setting each byte to zero. If each byte is set to zero then all the sprites will have a location of 0,0 and will all still be on screen when not active in the top left hand corner of the display.

It is far better to set the Y/X coords to 192,0 respectively which means that any non active sprites will be hidden below the bottom border.

The hardware of the machine will service all sprites within the sprite attribute table, including those which you may not be using. It is sometimes desirable to 'lock off' part of the sprite attribute. This can be done by writing a 208 (Decimal) into the Y position of the next sprite after the last sprite you wish to use.

If for example:

1. We were using 8 times 7 bit sprites
2. The first two shapes in the sprite generator table were a diamond and a square in order
3. The sprite attribute table was located at £1C00 (Hex), 7K
4. We were using the VRAM I/O routines described in section 1

Then the routines to perform the following tasks are listed below.

1. Clear the sprite attribute table
2. Lock off all but the first two sprites within the sprite attribute table
3. Display :
 - a) A yellow diamond sprite in the top left hand corner of the screen
 - b) A blue circle sprite in the top right hand corner of the screen

```
;
;
; SPRITE – NO REGISTER SET-UP REQUIRED ON ETRY
;           AF,BC,DE AND HL REGISTERS AFFECTED ON EDIT
;
;
;           USES TABLE VDTL (YELLOW DIAMOND TOP LEFT)
;           BCTR (BLUE CIRCLE TOP RIGHT)
;
;
;           VARIABLES SATSAD
;
;
SPRITE:    CALL    CLRSPR    ;Clear sprite attribute table
           LD      A,2      ;Lock off all but the first
                           ;two sprites within the sprite
           CALL    BLKSPR    ;attribute table
```

```

CALL      DRWSPR      ;Display two sprites on-screen
RET      ;Return to calling routine
;
;CLRSPR-CLEAR SPRITE ATTRIBUTE TABLE
;
;      NO REGISTER SET UP REQUIRED ON ENTRY
;
;      VRAM START ADDRESS OF SPRITE ATTRIBUTE TABLE
;      ASSUMED TO BE HELD IN TWO BYTE LOCATION SATSAD
;      (SPRITE ATTRIBUTE TABLE START ADDRESS)
;
SATSAD:  DW      0      ;SATSAD is a two byte location
;              ;which MUST hold the start
;              ;address of the sprite attribute table
;
CLRSPR:  LD      DE,(SATSAD) ;Set write to VRAM pointer to
CALL     VSETOT ;start of sprite attribute table
LD      B,32      ;Set loop counter = 32
CLRSP1:  LD      C,192     ;Set first byte to be output
;              ;(sprite Y coord) to 192
CALL     VDOUTP    ;and output byte to sprite
;              ;attribute table
LD      C,0       ;Set output byte to zero
CALL     VDOUTP    ;and output it 3 times to
CALL     VDOUTP    ;the remaining three bytes
CALL     VDOPUTP   ;within each control block of the
;              ;sprite attribute table
DYNZ    CLRSP1    ;Decrement loop counter and
;              ;if loop counter < > 0 then
;              ;goto CLRSP1 else drop through
;              ;to CLRSP2
CLRSP2:  RET      ;Return to calling routine
;
;BLKSPR-LIMIT NUMBER OF ACTIVE SPRITES TO VALUE CONTAINED IN
;      A ON ENTRY
;
;      IF A = 0 THEN NO ACTIVE SPRITES
;      IF A = (N) THEN (N) ACTIVE SPRITES WHERE
;      1 <= N <= 31
;      IF A > 31 THEN REUTRN FROM ROUTINE AND NO ACTION
;      TAKEN
;
BLKSPR:  CP      32      ;If A > 31 on entry to THIS
;              ;routine then return to
RET      NC      ;calling routine else drop
;              ;through to BLKSP1
BLKSP1:  SLA     A       ;Multiply contents of A
SLA     A       ;by 4
LD      E,A     ;Add the value of A * 4 to
LD      D,0     ;the start address of the
LD      HL,(SATSAD) ;sprite attribute table
ADD     HL,DE

```

```

EX      DE,HL      ;Set the write to VRAM pointer
CALL   VSETOT    ;to give the correct VRAM
                          ;address of the control block
                          ;for the sprite which is to
                          ;hold the locking off byte
LD      C,208    ;Output the locking off byte
CALL   VDOUTP    ;to VRAM
RET     ;Return to the calling routine
;
;DRWSPR-DISPLAY TWO SPRITES ON-SCREEN
;
;      A YELLOW DIAMOND IN THE TOP LEFT HAND CORNER OF
;      THE SCREEN
;      A BLUE CIRCLE IN THE TOP RIGHT HAND CORNER OF THE
;      SCREEN
;
;      USES TABLE YDTL (YELLOW DIAMOND TOP LEFT)
;      BCTR (BLUE CIRCLE TOP RIGHT)
;
YDTL:  DB      1,1,0,11      ;Sprite 0 control information
                          ;Byte 1 – Set Y coord to 1
                          ;Byte 2 – Set X coord to 1
                          ;Byte 3 – Display pattern 0
                          ;Byte 4 – Pattern colour is dark
                          ;yellow
BCTR:  DB      1,247,1,5    ;Sprite 1 control information
                          ;Byte 1 – Set Y coord to 1
                          ;Byte 2 – Set X coord to 247
                          ;Byte 3 – Display pattern 1
                          ;Byte 4 – Pattern colour is light
                          ;blue
;
DRWSPR: LD      DE,(SATSAD)  ;Set write to VRAM pointer to
CALL   VSETOT    ;start of sprite attribute table
LD      B,8      ;Loop counter = 8
LD      HL,YDTL  ;Set pointer to start of sprite
                          ;control data
DRWSP1: LD      C,(HL)      ;Output the control byte
CALL   VDOUTP    ;pointed to by HL to VRAM
INC    HL        ;Increment pointer to look at next
                          ;data byte
      DJNZ    DRWSP1      ;Increment loop counter and if
                          ;loop counter < > 0 then goto
                          ;DRWSP1 else drop through to
                          ;DRWSP2
DRWSP2: RET     ;Return to calling routine

```

^(Illus 7 – Screen display of above)

Sprite Generator Table 4.3

The sprite generator table holds a library of potential sprite patterns. It is a maximum of 2048 bytes long and starts on a 2K boundary (see section 2.2 Register 6). It is split into 256 blocks of 8 bytes each.

When 8 times 8 bit sprites are being used there are 256 different possible patterns available for use at any one time. If however, 16 times 16 sprites are in use, each sprite pattern will take four 8 times 8 bit blocks to make a complete shape, and therefore only 64 patterns are available for use.

An important point to realise from this is that when you are using 8 times 8 bit sprites you can access them sequentially by displaying pattern numbers 0,1,2,3,4.....253,254,255 etc. But when accessing 16 bit by 16 bit sprites you need to count in blocks of four so that the pattern numbers will be something like 0,4,8,12,16.....244,248,252 etc.

The diagram below shows how 16 times 16 bit sprite patterns held in the sprite generator table are mapped onto the screen.

Screen display.

^(Illus 8 – Quadrant diagram)

Start of Sprite generator table → Bytes 0,1,2,3,4,5,6,7
= Quadrant A
Bytes 8,9,10,11,12,13,14,15
= Quadrant B
Bytes 16,17,18,19,20,21,22,23
= Quadrant C
Bytes 24,25,26,27,28,29,30,31
= Quadrant D

The most important point is that patterns within this table are assigned a pattern number on the basis of position. The actual manipulation of the sprite patterns has been discussed in the previous section 4.2.

Animating Sprites (Assembler) 4.4

As has been discussed in section 4.1, animating sprites on the MTX is a matter of changing one byte within the sprite attribute table to display a sequence of patterns already resident in VRAM and held in the sprite pattern generator table.

The most difficult part of animating sprites is to get the timing right between displaying each sequential pattern. The best solution, which unfortunately is the most complex, is to perform this operation under interrupts. This has been separately described in section 3.

An alternate solution which works quite well is to use a descending loop counter and then display a new pattern in the animation sequence each time the counter reaches zero. The pattern numbers to be displayed are held within a table marked by a delimiter such as £FF (Hex), 255 (Decimal). This will allow the animation routine to detect the end of the animation sequence and wrap around to the start of the table to perform the sequence again without constraining it to a fixed number of display patterns between the start of the sequence and the end.

Shown below is a very simple sequence for a growing square consisting of four animation stages.

^(Illus 9 – Diagram of simple animation sequence).

The assembler code which will perform this animation is listed below.

The assumptions made are:

1. That VDP write only register 1 has been set up to enable 8 times 8 bit sprites with zero magnification (See section 2.2 register 1).
2. That these patterns have already been loaded into the sprite generator table into pattern positions 0,1,2 and 3.
3. That the sprite display has already been taken care of by other routes (See the previous section 'Setting up the sprite attribute table')
4. That the sprite attribute table is located at VRAM address £1C00 (Hex).
5. That we are using the first sprite control block in the sprite attribute table and therefore displaying sprite 0.
6. The routines will be making use of the VDP I/O routines described in section 1.1

This animation sequence can be altered to give an oscillating square with ease without affecting the routine in any way merely by changing the contents of the table ASTFS0 from 0,1,2,3,255 to 0,1,2,3,2,1,255

^(Illus 9a – Diagram of extended simple animation seq.)

```
;  
;  
;ANIMAT- ANIMATE SPRITE 0 – SEQUENCE IS GROWING SQUARE  
;      USES SPRITE PATTERNS 0,1,2 AND 3  
;      8 TIMES 8 BIT SPRITE  
;  
;  
;      NO REGISTER SET UP REQUIRED ON ENTRY  
;  
;  
;      NO REGISTERS AFFECTED ON EXIT  
;  
;  
;      THE ROUTINE MAKES USE OF A DESCENDING LOOP COUNTER  
;      CALLED ALCFS0  
;      ANIMATION LOOP COUNTER FOR SPRITE 0  
;  
;
```


CALL	VSETOT	;address which contains sprite
LD	C,A	;0 pattern number
CALL	VDOUTP	;Write pattern number to VRAM
RET		;updating previous pattern
		;Return to calling routine

This routine may animate a sprite too quickly or perhaps too slowly, in which case the speed of animation can be altered by changing the values contained in ALCREV and ALCFS0. A larger value will give a slower speed and a smaller value will give a faster speed. ALCREV and ALCFS0 must always be greater than one prior to calling the routine.

Sprite Impact Detection Routine (Assembler)

The VDP chip contains a sprite coincidence flag. It is bit 6 in the read only register (See Appendix B). Each time any two sprites coincide this bit is set, but it is not really much use because:

1. It does not tell you which two sprites have coincided.
2. It may not necessarily detect on the sprite shape itself because it senses on all of the pixels within the whole sprite block. Even those which are transparent.
3. It detects sprites which are not actually on the display screen if the X/Y coord values will cause overlap.

What is needed is a routine that will solve all of these problems, and this is described below.

The routine makes some assumptions:

1. That any non displayed sprites have X/Y coords of 0,192 respectively (see section 4.2 'Setting up the sprite attribute table')
2. That the routine is to check for impact relative to a particular sprite. In this case sprite 0, the first sprite within the sprite attribute table.
3. The routine will make use of the VDP I/O routines described in section 1.1.

An important point to note is that this routine can be used equally well with 8 times 8 bit sprites and 16 times 16 bit sprites. The value of OFFSET in the routine COMTSB is a key variable.

^(Illus 10 – Diagram of OFFSET range – SIDR)

As can be seen from the diagram above OFFSET narrows down the pixel area under which sprite impact is said to have occurred relative to the top left hand corner of the sprite. By using larger values of OFFSET it is possible to provide proximity testing prior to sprite impact because the routines below will then be looking at an area which is larger than the actual area of the sprite pattern.

With an OFFSET value of 4 the routines below are designed to detect impact between 8 and 9 bit sprites within a central area of 4 pixels

```

;
; IMPACT-TEST FOR IMPACT BETWEEN SPRITE 0
;       AND ANY OTHER ON SCREEN SPRITES
;
;       THE VRAM START ADDRESS OF THE SPRITE
;       ATTRIBUTE TABLE IS HELD IN THE TWO BYTE LOCATION
;       SATSAD (SPRITE ATTRIBUTE TABLE START ADDRESS)
;       -       -       -       -       -
;
;       NO REGISTER SET UP REQUIRED ON ENTRY
;       NO REGISTERS AFFECTED ON EXIT
;
;       IF IMPACT TRUE THEN VARIABLE TRUFAL = 1
;       ELSE TRUFAL = 0
;
;       IF TRUFAL = 1 THEN 7 BYTE TABLE IMSPR CONTAINS
;
;       1) OFFENDING SPRITE NUMBER           - BYTE 1
;       2) OFFENDING SPRITE VRAM ADDRESS     - BYTE 2,3
;       3) OFFENDING SPRITE X COORD         - BYTE 4
;       4) OFFENDING SPRITE Y COORD         - BYTE 5
;       5) OFFENDING SPRITE PATTERN NUMBER  - BYTE 6
;       6) OFFENDING SPRITE COLOUR          - BYTE 7
;
SATSAD:  DW      £1000           ;See description above
TRUFAL:  DB      0              ;See description above
;
IMSPR:   DB      0              ;See description above
         DB      0,0           ;See description above
         DB      0              ;See description above
;
IMPACT:  PUSH     AF            ;Save old Acc and flags
         PUSH     BC            ;Save old BC register pair
         PUSH     DE            ;Save old DE register pair
         PUSH     HL            ;Save old HL register pair
         PUSH     IX            ;Save old IX register pair
;
         LD      DE,(SATSAD)    ;Set read from VRAM pointer
         CALL    VSETRD         ;to start of sprite attribute table
         LD      A,0            ;Set impact true/false byte
         LD      (TRUFAL),A     ;to false
         CALL    VDINPT         ;Read sprite 0 X/Y coords
         LD      E,C            ;from VRAM sprite attribute
         CALL    VDINPT         ;table into DE register pair
         LD      D,C            ;Register D = X coord
         LD      E,C            ;Register E = Y coord

```

```

LD      A,E      ;Check to see that sprite o
CP      192      ;is actually on screen
JR      NZ,IMPAC0 ;and if condition true goto
LD      A,D      ;IMPAC0
CP      0        ;else return to calling
RET     Z        ;routine with variable
;TRUFAL = false

;
IMPAC0: CALL     VDINPT ;Move read from VRAM pointer
CALL     VDINPT ;onto next sprite
LD      B,31     ;Set loop counter = 31

;
IMPAC1: PUSH    BC    ;Save loop counter
CALL     VDINPT ;Read sprite Y coord from
;VRAM sprite attribute table
LD      A,C      ;and test to see whether the
CP      208      ;routine has detected a user
;defined end of sprite
JR      NZ,IMPAC2 ;attribute table marker
;If condition is false then
POP     BC       ;goto IMPAC2
;else retrieve old loop
JP      IMPAC4   ;counter value and goto
;IMPAC4 to exit routine

;
IMPAC2: LD      L,A    ;Read sprite X/Y coords from
CALL     VDINPT ;VRAM sprite attribute table
LD      H,C      ;into HL register pair
;Register H = X coord
;Register L = Y coord
CALL     COMTSB ;Call impact detection test
;routine
LD      A,(TRUFAL) ;and if the variable TRUFAL
CP      0        ;equals 0 on exit then sprite
JP      NZ, IMPAC3 ;impact has not occurred
;therefore continue testing
;routine and move onto next
;sprite else goto IMPAC3
; and set up table IMPSPR
CALL     VDINPT ;Move read from VRAM pointer
CALL     VDINPT ;onto next sprite
POP     BC       ;Retrieve old loop counter
;value and decrement it
DJNZ    IMPAC1  ;If loop counter < > 0 then
JP      IMPAC4  ;goto IMPAC1 and repeat
;routine else goto IMPAC4
;and exit routine

;
IMPAC3: POP     BC    ;Retrieve old loop counter
;value
LD      IX,IMPSPR ;Set index register pointer
;to point to start of table IMPSPR
LD      (IX+3),H ;Byte 4 of table IMPSPR

```

```

LD      (IX+4),L      ;= offending sprite X coord
                        ;Byte 5 of table IMPSPR
CALL    VDINPT        ;= offending sprite Y coord
LD      (IX+5),C      ;Byte 6 of table IMPSPR
CALL    VDINPT        ;= offending sprite pattern number
LD      (IX+6),C      ;Byte 7 of table IMPSPR
LD      A,32          ;= offending sprite pattern colour
LD      B              ;Byte 1 of table IMPSPR
SUB     B              ;= offending sprite number
LD      (IX+0),A
SLA    A              ;Determine actual VRAM
SLA    A              ;address of offending sprite
LD      E,A           ;using the reference value
LD      D,0           ;start address for the VRAM
LD      HL,(SATSAD)   ;sprite attribute table held
ADD    HL,DE          ;in the two byte location
                        ;SATSAD and the loop counter
                        ;value held in register B
LD      (IX+1),L      ;Bytes 2 and 3 of table
LD      (IX+2),H      ;IMPSPR hold two byte VRAM
                        ;address of offending sprite
;
IMPAC4: POP    IX      ;Reset old IX register pair
        POP    HL      ;Reset old HL register pair
        POP    DE      ;Reset old DE register pair
        POP    BC      ;Reset old BC register pair
        POP    AF      ;Reset old Acc and flags
        RET           ;Return to calling routine
;
; COMTSB-COMPARE TWO SPRITE BLOCKS
; --- - - - -
;
;
; COMPARES TWO SPRITE BLOCKS
; ONE SET X,Y IN HL RESPECTIVELY
; THE OTHER, X,Y IN DE RESPECTIVELY
; AND RETURNS TRUFAL = 1 IF TRUE
; ELSE TRUFAL = 0
;
;
; SPRITE X COORDS MUST NOT EQUAL 0 AND
; SPRITE Y COORDS MUST NOT EQUAL 192
;
OFFSET: DB      0      ;Reserve 1 byte for variable
                        ;OFFSET
;
COMTSB: LD      A,4    ;Set variable OFFSET to 4
        LD      (OFFSET),A
        LD      A,0    ;Set impact true/false
        LD      (TRUFAL),A ;variable to false
;

```

```

COMTS1:  LD      A,L      ;Check to see that sprite
         CP      192     ;is on-screen
         JR      NZ,COMTS2 ;If condition true then goto
         LD      A,H     ;COMTS2 else return to
         CP      0      ;routine with TRUFAL = false
         RET     Z
;
COMTS2:  LD      A,(OFFSET) ;Perform impact true/false
         ADD     A,L     ;test on each of the four
         CP      E      ;sides of the two active
         JR      C,COMTS3 ;sprites
         LD      A,(OFFSET) ;If at any point the test
         ADD     A,E     ;routine shows that the
         CP      L      ;two sprites are not close
         JR      C,COMTS3 ;enough for impact to take
         LD      A,(OFFSET) ;place then an exit is made
         ADD     A,H     ;to COMTS3 and TRUFAL is
         CP      D      ;false
         JR      C,COMTS3 ;Only if ALL impact test
         LD      A,(OFFSET) ;conditions are true will
         ADD     A,D     ;TRUFAL be set to true
         CP      H      ;when the test routine
         JR      C,COMTS3 ;is completed
         LD      A,1
         LD      (TRUFAL),A
         RET
;
COMTS3:  NOP
         RET
;

```

Looking underneath sprites at the background

In many situations it may be desirable or even necessary to look underneath sprites at objects or regions contained within the pattern background.

There is a direct and simple proportional relationship between the pattern background and the position of sprites on screen.

Referring to scion 2.4 and assuming for the moment that this routine is being written for a graphics 2 mode application, and that we are using 16 times 16 bit sprites, there are actually two ways in which this problem can be tackled.

If we re using a high resolution plottable type graphics 2 mode set up (as BASIC does), then the routine will not detect characters underneath sprites, but will detect bits within pattern generator table bytes and will act on the pattern generator table. This code is listed in Option 1.

If we are using the other option which is to set up three identical pattern libraries within the pattern generator table and to have a dynamically changing screen (Pattern name table), the routine will operate on the pattern name table and will return characters as results. This section of code is listed in Option 2.

The relationship between both of these options and the sprites are shown in the diagram below.

^(Illus 11 – Option 1 sprite to background mapping)

^(Illus 12 – Option 2 sprite to background mapping)

Each code option makes use of offset values to test underneath points within the sprite relative to the top left hand corner of the sprite. Simply by changing the values of these offsets it is possible to detect on new areas of the sprite. Also by adding or deleting new calls within the DOT routine it is possible to detect on more or less points within any sprite.

Both code sections listed in options 1 and 2 make use of the 1 byte variables SPRP1 to SPRP12, and the 2 byte variables CENLTL, CENLTR, CENLBL and CENLBR.

As shown in the diagram below SPRP1 to SPRP12 are actually points within a sprite pattern area, and will contain the contents of the pattern name or generator table bytes which map underneath them.

The variables CENLTL, CENLTR, CENLBL and CENLBR, will contain the pattern name or generator table addresses of the bytes held in SPRP5, SPRP8 and SPRP9 respectively.

^(Illus 13 – Diagram of sprite points on sprite)

In the case of option 1, DOT will set SPRP1 to SPRP12 to 1,s or 0,s dependent on whether a bit is set or no set underneath the sprite at the corresponding point on screen.

In the case of option 2, DOT will set SPRP1 to SPRP12 to the character value underneath the sprite at the corresponding point on screen.

Option 1

The code for option 1 is identical to the code for option 2 with two major differences.

1. The routine DOTSUB is different and is listed below. It does not return character values in SPRP1 to SPRP12 but returns bit set or not set indicated by either 1 or 0. It does not operate on the pattern name table but uses the pattern generator table, the start address of which is held in a two byte location SAPATG.

2. There is an additional table BITTAB used by DOTSUB.

```

;
; DOTSUB – DERIVES HIGH-RES SCREEN LOCATION OF A POINT WITHIN
; A SPRITE PATTERN AND TEST BIT UNDERNEATH THAT POINT
; RETURN RESULT OF TEST IN ACC
; IF BIT SET THEN ACC = 1
; IF BIT NOT SET THEN ACC = 0
;
; X COORD IN H, Y COORD IN L ON ENTRY
;
; PATTERN GENERATOR ADDRESS IN DE
; BYTE IN A
; ON EXIT
;
; AF, BC, DE, HL REGISTER PAIRS AFFECTED ON EXIT
;
; ROUTINE USES TWO BYTE VARIABLE SAPATG WHICH MUST
;
; BE SET TO CONTAIN THE START ADDRESS OF THE PATTERN
; GENERATOR TABLE PRIOR TO ENTRY TO THIS ROUTINE
;
; BITTAB IS AN 8 BYTE TABLE USED TO PERFORM A BIT TEST
; ON INDIVIDUAL BYTES EXTRACTED FROM THE PATTERN
; GENERATOR TABLE
;
SAPATG:  DW      £0000          ;(See description above)
BITTAB:  DW      128,64,32,16  ;(See description above)
         DB      8,4,2,1,0
;
DOTSUB:  PUSH     HL           ;Save sprite X/Y coords
         LD      A,L         ;Derive relative Y offset from
         LD      D,0         ;start of pattern generator table
         AND     248         ;Result in DE when this section
         LD      E,A         ;finished
         LD      B,5         ;DE = INT ( Y / 8 ) * 256 +
DOTSUB:  SLA     E           ; ( Y – (INT ( Y / 8 ) * 8 ))
         RL      D
         DJNZ   DOTSU1
         LD      A,L
         AND     7
         AND     A,E
         LD      E,A         ;Y coord offset in DE at this point
;
         LD      A,H         ;Derive relative X offset from
         LD      H,0         ;start of pattern generator table
         AND     7           ;Result in HL when this section
         LD      L,A         ;finished
         ADD     HL,DE       ;HL = INT ( X / 8 ) * 8
         LD      DE, (SAPATG) ;Add Y offset to X offset
         ;Result in HL
         ;Actual start address of pattern
         ;generator table to DE

```

ADD	HL, DE	;Add to relative address contained
		;in HL to give true VRAM address
EX	DE, HL	;of pattern generator byte to be
		;examined – Result in DE
POP	HL	;Retrieve old sprite X/Y coords
PUSH	DE	;Save pattern generator byte
		;address
LD	A,H	;Perform bit test on extracted
AND	248	;pattern generator byte
LD	L,A	;using table BITTAB
LD	H,0	
LD	DE,BITTAB	
ADD	HL,DE	
LD	A,(HL)	;Read byte from pattern generator
POP	DE	;table
CALL	VSETRD	
CALL	VDINPT	
AND	C	;If bit = 0 then exit routine
CP	0	;Acc = 0
RET	Z	
LD	A,1	;If bit is < > then Acc = 1
RET		;Return to calling routine

Option 2

```

;
;
;DOT -SPRITE TO BACKGROUND MAPPING ROUTINE
;
;
; LOOK AT 12 POINTS WITHIN 16 6IMES 16 BIT SPRITE PATTERN
; AND RETURN THE CHARACTER VALUE UNDERNEATH THE
; SPRITE EXTRACTED FROM THE PATTERN NAME TABLE
; (DISPLAY SCREEN) AT EACH OF THESE POINTS IN SPRP1
; TO SPRP12
;
;
; RETURN THE PATTERN NAME TABLE (DISPLAY SCREEN)
; ADDRESSES OF THE FOUR CENTRE POINTS OF THE SPRITE
; PATTERN WITHIN THE LOCATIONS:
;
;
; CENLTL - CENTRE LOCATION TOP LEFT QUADRANT
; CENLTR - CENTRE LOCATION TOP RIGHT QUADRANT
; CENLBL - CENTRE LOCATION BOTTOM LEFT QUADRANT
; CENLBR - CENTRE LOCATION BOTTOM RIGHT QUADRANT
;
;
; REGISTER PAIR DE MUST CONTAIN THE VRAM START ADDRESS
; OF THE CONTROL BLOCK FOR THE SPRITE TO BE TESTED ON
; ENTRY
;
;
; THE SPRITE TO BE TESTED MUST BE ON-SCREEN WHEN DOT
; IS CALLED
;
;
;

```



```

;
DOT4:   LD      A,L      ;Increment Y coord in register L
        ADD     A,7     ;by 7 to move onto right middle
        LD      L,A     ;side of sprite
        PUSH   HL      ;Save HL register pair
        CALL   DOTSUB  ;(See description in section DOT1)
        POP    HL      ;Retrieve saved HL register pair
        LD     (SPRP7),A ;Store character value under right
                        ;middle side of sprite in SPRP7

;
DOT5:   LD      A,H     ;Decrement X coord in register H
        SUB     7      ;by 7 to move onto the top right hand
        LD     H,A     ;quadrant centre of sprite
        PUSH   HL      ;Save HL register pair
        CALL   DOTSUB  ;(See description in section DOT1)
        POP    HL      ;Retrieve saved HL register pair
        LD     (SPRP6),A ;Store character value under top right
                        ;hand quadrant centre of sprite in SPRP6

;
DOT6:   DEC     H       ;Decrement X coord in register H
                        ;by 1 to move onto top left hand
                        ;quadrant centre of spite
        PUSH   HL      ;Save HL register pair
        CALL   DOTSUB  ;(See description in section DOT1)
        POP    HL      ;Retrieve saved HL register pair
        LD     (SPRP5),A ;Store character value under top left
                        ;hand quadrant centre of sprite in SPRP5

;
DOT7:   LD      A,H     ;Decrement X coord in register H
        SUB     7      ;by 7 to move onto left middle
        LD     H,A     ;side of sprite
        PUSH   HL      ;Save HL register pair
        CALL   DOTSUB  ;(See description in section DOT1)
        POP    HL      ;Retrieve saved HL register pair
        LD     (SPRP4),A ;Store character value under left
                        ;middle side of sprite in SPRP4

;
DOT8:   INC     L       ;Decrement Y coord in register L by 1
        LD     A,H     ;and increment X coord in register H
        ADD     A,7     ;by 7 to move onto bottom left hand
        LD     H,A     ;quadrant centre of sprite
        PUSH   HL      ;Save HL register pair
        CALL   DOTSUB  ;(See description in section DOT1)
        POP    HL      ;Retrieve saved HL register pair
        LD     (SPRP8),A ;Store character value under bottom
                        ;left hand quadrant centre of sprite
                        ;in SPRP8
        LD     (CENLBL),DE ;Store address of character value held
;in SPRP8 in two byte location CENLBL
;

```

```

DOT9:      INC      H          ;Increment X coord held in register H
          ;by 1 to move onto bottom right hand
          ;quadrant centre of sprite
          PUSH     HL         ;Save HL register pair
          CALL     DOTSUB    ;(See description in section DOT1)
          POP      HL         ;Retrieve saved HL register pair
          LD       (SPRP9),A ;Store character value under bottom
          ;right hand quadrant centre of sprite
          ;in SPRP9
          LD       (CENLBR),A ;Store character value under bottom
          ;right hand quadrant centre of sprite
          ;in SPRP9

```

```

;
DOT10:    DEC      H          ;Decrement X coord held in register H
          LD       A,L        ;by 1 and increment Y coord held in
          ADD      A,7        ;register L by 7 to move onto bottom
          LD       L,A        ;middle of sprite
          PUSH     HL         ;Save HL register pair
          CALL     DOTSUB    ;(See description in section DOT1)
          POP      HL         ;Retrieve saved HL register pair
          LD       (SPRP11),A ;Store character value under bottom
          ;middle of sprite in SPRP11

```

```

;
DOT11:    LD       A,H        ;Decrement X coord held in register H
          SUB      7          ;by 7 to move onto bottom left hand
          LD       H,A        ;corner of sprite
          PUSH     HL         ;Save HL register pair
          CALL     DOTSUB    ;(See description in section DOT1)
          POP      HL         ;Retrieve saved HL register pair
          LD       (SPRP10),A ;Store character value under bottom
          ;left hand corner of sprite in SPRP10

```

```

;
DOT12:    LD       A,H        ;Increment X coord held in register H
          ADD      A,15       ;by 15 to move onto bottom right hand
          LD       H,A        ;corner of sprite
          PUSH     HL         ;Save HL register pair
          CALL     DOTSUB    ;(See description in section DOT1)
          POP      HL         ;Retrieve saved HL register pair
          LD       (SPRP12),A ;Store character value under bottom
          ;right hand corner of sprite in SPRP12
          RET          ;Return to calling routine

```

```

;
;DOTSUB-DERIVES SCREEN LOCATION OF A POINT WITHIN A SPRITE
PATTERN AND BYTE HELD WITHIN THE SCREEN ADDRESS

```

```

;
;      X COORD IN H, Y COORD IN L ON ENTRY

```

```

;
;      LOCATION ON SCREEN IN DE
;      BYTE IN A
;      ON EXIT

```

```

;
;      AF, BC, DE, HL REGISTER PAIRS AFFECTED ON EXIT

```


CALL	VDINPT	;Read sprite X coord from VRAM and
LD	H,C	;place into register H
POP	BC	;Retrieve old BC Register pair
RET		;Return to calling routine

Joystick and keyboard control – Introduction

5.1

On the MTX series computers, the joystick (left hand side), is mapped onto the cursor keys and home key as shown below.

Cursor left	Joystick left
Cursor right	Joystick right
Cursor up	Joystick up
Cursor down	Joystick down
Home key	Fire key

What this effectively means is that if you wish to manipulate the joystick ports and read joystick data, you can dispense with having to look at the joystick ports and simply scan the keyboard for the corresponding cursor results.

This is a very simple process to perform but it does have what may be a disadvantage in some applications in that it cannot detect multiple key presses at the same time.

However, this can be done by performing a strobe across selected lines of the keyboard, and then reading the results directly from the read lines as described in sections 5.2 and 5.3 by by-passing the operating systems. 'GET A CHARACTER' routine.

The sense lines and read lines for the keyboard on the MTX are tied into the same port. This is port 5.

An output to this port is interpreted as a sense byte
An input from this port is interpreted as a read byte

An important point to realise is that the sense lines and read lines are active on bit low. That is key presses across the keyboard matrix are indicated by 0,s rather than 1,s.

The difficult part of writing a routine to perform keyboard scanning is selecting the right values to output along the sense lines of the keyboard, and knowing which lines to look at when examining the value returned from the read lines.

In both code options described in section 5.2 and 5.3 the correct scan values have already been worked out for the cursor keys. If however, you wish to create your won keyboard scan routines to look at other keys, the correct values can be obtained by reading appendix D (keyboard layout), and appendix E (keyboard scan values).

Joystick Manipulation (BASIC)

5.2

This section describes the basic code which will perform a keyboard strobe on cursor left/right/up/down), and home.

```
100 LET SENSEBYTE=223 : OUT 5,SENSEBYTE
    : LET READBYTE =INP(5)
110 IF READBYTE=127 THEN (firekey has been pressed)
120 LET SENSEBYTE=247 : OUT 5,SENSEBYTE
    : LET READBYTE=INP(5)
130 IF READBYTE=127 THEN (cursor left has been pressed)
140 LET SENSEBYTE=129 : OUT 5,SENESBYTE
    :LET READBYTE=INP(5)
150 IF READBYTE=127 THEN (cursor right has been pressed)
160 LET SENSEBYTE=251 : OUT 5,SENSEBYTE
    : LET READBYTE=INP(5)
170 IF READBYTE=127 THEN (cursor up has been pressed)
180 LET SENSEBYTE=191 : OUT 5,SENSEBYTE
    : LET READBYTE=INP(5)
190 IF READBYTE=127 THEN (cursor down has been pressed)
```

Joystick Manipulation (Assembler)

5.3

This section describes the machine code which performs the same task as the BASIC described above.

The routine is called GETCHR and is relocatable. It requires no parameters on entry, and will return the results of the keyboard strobe in a variable called KEYCHR according to the table listed below.

No key pressed	KEYCHR = 0
Cursor left	Bit 0 set
Cursor right	Bit 1 set
Cursor up	Bit 2 set
Cursor down	Bit 3 set
Home key	Bit 4 set

This table assumes that bit 0 is the LSB, and bit 7 is the MSB.

```
;
;
; GETCHR-USES VARIABLE 'KEYCHR'
; NO PARAMETERS ON ENTRY REQUIRED
; AF AND HL REGISTERS AFFECTED ON EXIT
;
; 'KEYCHR' SET ACCORDING TO TABLE ABOVE
;
;
KEYCHR:  DB      0          ;See description above
;
```

```

GETHCHR: LD      A,0
          LD      HL(KEYCHR) ;Set KEYCHR to zero – no key
          LD      (HL),A      ;pressed
;
          LD      A,223      ;Select strobe byte to scan
          OUT     (5+OFFSET),A ;fire key and output to port 5
          IN      A,(5+OFFSET) ;Examine read lines and if fire
          CP      127        ;key is not depressed then goto
          JR      NZ,GETLEF   ;GETLEF else
          SET     4,(HL)      ;Site fire key bit (condition true)
;
GETLEF:  LD      A,247      ;Select strobe byte to scan
          OUT     (5+OFFSET),A ;cursor left key and output to
          IN      A,(5+OFFSET) ;port 5 – examine read lines and
          CP      127        ;if cursor left is not depressed
          JR      NZ,GETRGT   ;then got GETRGT else
          SET     0,(HL)      ;Set cursor left bit (condition true)
;
GETRGT:  LD      A,239      ;Select strobe byte to scan
          OUT     (5+OFFSET),A ;cursor right key and output to
          IN      A,(5+OFFSET) ;port 5 – examine read lines and
          CP      127        ;if cursor right is not depressed
          JR      NZ,GETUP    ;then got GETUP else
          SET     1,(HL)      ;Set cursor right bit (condition true)
;
GETUP:   LD      A,251      ;Select strobe byte to scan
          OUT     (5+OFFSET),A ;cursor up key and output to
          IN      A,(5+OFFSET) ;port 5 – examine read lines and
          CP      127        ;if cursor up is not depressed
          JR      Z,GETUP     ;then got GETDWN else
          SET     2,(HL)      ;Set cursor up bit (condition true)
;
GETDWN:  LD      A,191      ;Select strobe byte to scan
          OUT     (5+OFFSET),A ;cursor down key and output to
          IN      A,(5+OFFSET) ;port 5 – examine read lines and
          CP      127        ;if cursor down is not depressed
          RET     NZ,        ; then return to calling routine
          SET     3,(HL)      ;else set cursor down bit
                                ;(condition true)
;
          RET                ;Return to calling routine

```

Get a character ROM routine.

5.4

In m/c code it is easy to get a character by using the call KBD option. The entry point for DBD is located at £0079. It is designed to be transparent to the user and does not affect any registers. The only bit affected is the zero flag.

The result is returned in A and will be a standard ASCII value or a specific non-ASCII MTX keyboard value.

If the zero flag is set then no key has been pressed

If the zero flag is not set then a key has been pressed

To access KBD perform the command 'CALL £0079'

An important point to note concerning KBD is that it makes use of the variable LASTKEY to remove debouncing problems.

KBD is normally used as part of BASIC and will therefore perform BASIC error and escape sequences. If you do not wish the break key to be serviced, then you must disable it prior to calling this routine through the location INTFFF £FD5E (Hex), 64826 (Decimal).

This location can also be used to enable or disable auto-repeat, sound, sprite movement and cursor flash.

Referring to the table below. If a bit is set then the feature it enables is ON. If a bit is not set then the feature it enables is OFF.

INTFFF	Bit 0	Sound
	Bit 1	Break key
	Bit 2	Keyboard auto repeat
	Bit 3	Sprite movement and cursor flash
	Bit 4	User 1
	Bit 5	User 2
	Bit 6	User 3
	Bit 7	(Unused at present)

Another important location is be aware of is KBDFLG at location £FA91 (Hex), 64145 (Decimal). This provides three more features listed below. If a bit is set then the feature it enables is ON. If a bit is not set then the feature it enables is OFF.

KBDFLG	Bit 7	Alpha lock
	Bit 5	Page/Scroll
	Bit 2	Numeric keypad

Points to Look Out For 5.5

If you are running a section of code and for some reason BASIC is disabled the KBD routine described in section 5.4 will not work. It may be necessary to run an edited version of KBD in RAM which is completely independent of BASIC as an integral part of you programme. This alternative version of KBD is listed in appendix F.

If you are considering writing an item of software for the MTX which is joystick based, it is worthwhile noting that the vast majority of software uses the right joystick as describe in this section (see guide to prospective software writers – section 10).

Data Output/Input To/From Tape – Introduction

6.1

In some applications it may be desirable or even necessary to save and load variables or blocks of memory independently of your programme in files or discrete segments. For most people who have an MTX at this time the cost of a disk system may be too much to afford, and therefore the only logical alternative is to perform data file manipulation on tape.

The MTX cassette system is relatively fast (2400 baud approx), and has proven to be very reliable in extensive field trials, so this alternative may not be too bad.

Essentially all data I/O is channelled through a routine called INOUT which is located at £0AAE (Hex). Most of the code written for this section will make use of this routine.

Tape I/O using ROM routines (Assembler)

6.2

The tape I/O ROM routine INOUT located at £0AAE (Hex) on its own, is designed to save and load blocks of raw data extracted directly from memory and only from the first BASIC page.

Variable retrieval and page control are taken care of by other routines and are discussed in later sections.

A memory block save/load routine requires three parameters to be set up prior to execution. These are listed below:

- i. Register pair HL must contain the start address at which data bytes are to be transferred from/to
- ii. Register pair DE must contain the number of bytes which are to be transferred to/from tape.
- iii. Variable TYPE located at £FD68 (Hex) must contain either :
 - a. 0 – To indicate 'Save Data'
 - b. 1 – To indicate 'Load Data'

A section of code which will save 100 (Decimal) bytes starting from location 10,000 (Decimal) is listed below.

There are several important points to note concerning INOUT:

- i. The routine INOUT services the break key during execution. If this option is not required then the break key must be disabled through the location INTFFF (see section 5.4)
- ii. It does not create any form of tape header label or display any user screen prompts such as :

PRESS PLAY AND RECORD ON TAPE
AND HIT ANY KEY WHEN READY

and will begin saving or loading data the moment it is executed. This means that any programmer who wishes to make use of this routine must write his or her own prompts.

```

;
;DATAIO – SEE DESCRIPTION ABOVE FOR ROUTINE SUMMARY
;
;
;      PARAMETERS REQUIRED ON ENTRY
;      SEE DESCRIPTION ABOVE
;
;
;      AF, BC, DE, HL REGISTERS AFFECTED ON EXIT
;
;
DATAIO:  LD      HL,10000      ;Start address at which data
;                          ;is to be saved from
;      LD      DE, 100      ;Number of bytes to be saved
;      LD      A,0          ;Set 'save/load' option byte
;      LD      (£FD68),A    ;Type to save
;      CALL   £0AEE        ;Save data bytes
;      RET                    ;Return to calling routine

```

An edited version of the tape routine INOUT is listed within appendix G. It is not dependent on BASIC and can be used in applications where BASIC is not required to be in use.

Adding Data Tape Save/Load Commands to BASIC

6.3

This routine has been designed so that it can be added to any BASIC program or written as the core to a new BASIC program and will add two new commands listed below.

- i. USER,L,<string\$> - Load data string
- ii. USER,S,<string\$> - Save data string

Once the code has been written into the BASIC program it must be executed before the new form of USER is typed in via the editor otherwise the command USER will be rejected.

it has been designed so that it will save and load strings only.

An important point to note is that all strings must have a length of some sort when using this routine. If we were saving a string A\$ which was equal to "" (Null), its apparent length would be zero but internally this would be treated as 64K. In order to give null strings a length they can be assigned with the value CHR\$(0) or a spce.

??? (Insert rest of data tape save/load section here).

References and Acknowledgements

7.8

References

1. Texas TMS9929L Technical Handbook
2. Zilog Z80 Handbook
3. Zilog CTC Handbook
4. Memotech customer enquiries

Acknowledgements

Many thanks to:

1. David Fazackerley
2. David Netherwood
3. L R Whalley
4. Jeff Wakeford (Artwork)

And in particular all of Memotechs customers for the large amount of interest shown in the MTX.

A guide for prospective software writers – Introduction

8.1

The following is intended to help software writers get their programs accepted, and published quickly by Continental Software. In no way does compliance with the following guarantee publication and should be considered only as a guide.

Documentation

8.2

All programs should be accompanied by full instructions, preferably typed. Games should come with details of the keys to be used, the object of the game, features of different levels of play, the score system and any other details which may be of concern to us or our customers.

Business software should include details of the relevant standards to which it conforms, eg type of accounting procedure, Country in which it is intended for use and possible applications, eg small business, chemists etc.

Program Standards

8.3

All programs must be crash proof, ie be able to cope with all possible combinations of input and user errors. Where the user does make an error the program should take some suitable course of action. In a games program this might be to ignore the erroneous input, in a business program this may mean displaying an error message and an invitation to re-enter the input.

Business and educational programs should have some form of validation to ensure that input is reasonable, eg range checks data type checks, format checks and any other reasonable precaution against entering invalid data by mistake.

Programs should on the whole be simple to use, where this is not possible they should include help pages. All programs should have instructions included in them.

Where joysticks or keys are to be used then we would prefer the use of the joystick marked right/cursor control keys (the two are mapped onto each other). However, we accept that it is sometimes more convenient to use the joystick marked left and its associated keys.

In general programs should be small enough to run on the 32k MTX 500 machine.

Presentation

8.4

Presentation is of paramount importance. It is often all a prospective purchaser has to go on, so please take some time and care over this very important aspect of programming.

Both games and business software should make good use of sound and colour where possible, though this should be done with care and taste. Where textual output is concerned it should be correctly spelt, properly punctuated, and well laid out. Numeric output should be tabulated where appropriate and column headings etc used. Where the option to print out is available facilities should be offered to re-direct output down the RS232 (see IOPR and IOPL system variables – Current BASIC manual).

Assembler Code Programs

8.5

Assembly code programs **MUST** contain all their labels, and all internal program references must be by label. This is so we can relocate the program to run on both the MTX 500 and 512. Please also note the following:

Labels must not be more than 6 characters in length

Program data lines should be kept to a reasonable length, ie less than 80 characters per line where possible.

Pure assembler code programs should not rely on the machine being in any particular state when it takes control, eg being in graphics mode etc. Our preference is for this type of program.

The hash character should not be used anywhere in the program except to denote hex numbers.

BASIC Programs

8.6

When writing business software and the like it is important to take great care over input and the ease of inputting data, as a rule BASICS own input facilities are not good enough

Blended BASIC and Assembler Code Programs

8.7

Where a combination of assembler and basic is used all the assembly code should be contained in one line at the beginning of the program. If you do not wish to execute this line at the start of the program make the first line of assembler a RET statement, rather than skip around the line altogether. Calls to machine code routines should then be made by means of the USR function.

NB it is important to mark the entry points used by the USR calls in the assembler code eg

10 code

```
4007          RET
4008 WAIT:    CALL £79 ; ENTRY POINT TO WAIT FOR KEY
400B          JR, Z,WAIT
400D          RET
```

SYMBOLS:

```
WAIT      4007
```

```
100 LET WAIT=4*1024+8
110 LET X=USR (WAIT)
120 STOP
```

Program Media

8.8

Programs can be accepted as cassettes, 40 track 5 ¼" floppy disks, provided they will run on an FDX (type 0-3 in config), 8" floppy disks config codes 10-13 or for CPM software LIFEBOAT FORMAT A1 (Type 10) disks.

Cassettes should have a recording on both sides and should clearly indicate which machine the program was written on.

Disks should likewise contain a backup copy of the program, together with any appropriate source code and details of any compilers, interpreters or assemblers used.

Technical/Commercial Enquiries 8.9

All technical enquiries should be directed to Technical Services at Memotech. Completed programs and commercial enquiries should be sent to:

Mr T Spencer
Software Co-ordinator
Memotech
Unit 23 Station Lane
Whitney
Oxon

Appendix A

7.1

VDP BASIC Memory Mapping

This section gives the actual start addresses of the various tables created in VRAM by BASIC to generate video displays and sprite patterns

Text Mode (VS 5)	Start Address
-------------------------	----------------------

Pattern name table (screen)	- 7K
Pattern generator table	- 6K

Graphics 2 Mode (VS 4)	Start Address
-------------------------------	----------------------

Pattern name table (screen)	- 15K
Pattern generator table	- 0 to 6143 (Decimal)
Pattern colour table	- 8K to 8K+6143
Sprite attribute table	- 15K+768
Sprite generator table	- 14K

Appendix B

7.2

Sprite Coincidence Flag

This is held in the VDP READ ONLY REGISTER shown in diagram 1.1 below.

(Illus 15 – Sprite coincidence register)

This register can be read and its contents examined from BASIC or machine code.

In BASIC `3000 LET READRG=INP(2)`

Will return contents of this location in variable READREG

In M/C CODE `IN A,(2)`

Will return contents of this location in A

The sprite coincidence flag, which detects impact of any two sprites will be 1 if impact has occurred, and 0 if it has not. However, the use of it requires caution because it will detect impact of zero sprites (that is the ones you are not using if the x/y coordinates match). You need some method of 'locking off' extra sprites, and this can be done by writing the decimal value 208 into the y coord position of the sprite number after the last sprite you wish to use.

In diagram 1.1 the sprite coincidence flag is bit 'C'. You can read this in either BASIC or M/C code as shown below.

In BASIC 3010 LET READREG=READREG AND 32
3020 IF READREG=1 THEN (condition true)
ELSE (false)

In M/C Code BIT (5), a
JP NZ, (condition true)
(false)

BASIC itself copies the contents of the VDP read only register into location £FE54 (Decimal ????), during the interrupt (Copy of VDP Status Register).

Appendix C

7.3

Colour Assignments

Colour Code Hex	Colour Code Dec	Colour
0	0	Transparent
1	1	Black
2	2	Medium green
3	3	Light green
4	4	Dark blue
5	5	Light blue
6	6	Dark red
7	7	Cyan
8	8	Medium red
9	9	Light red
A	10	Dark yellow
B	11	Light yellow
C	12	Dark green
D	13	Magenta
E	14	Grey
F	15	White

Appendix D

7.4

Keyboard Layout

The diagram below is a logical map showing how the keyboard is laid out from the machine hardware point of view.

Outputting the correct sense value via port 5 into the keyboard matrix and reading the correct value from port 5 as a scan byte is discussed in more detail in section 5.

A table summary of the values is given in Appendix E.

Lower case values are detected by looking at the shift key.

(Illus 14 – Diagram of keyboard layout)

Appendix E

7.5

Table Summary of Keyboard Sense/Scan Values

To use this table please refer to section 5 and Appendix D.

The table summary does not cover all combinations of keys but only:

1. Alphabetic A to Z
2. Digits 0 to 9
3. Cursor keys and home key

Using this table and the diagram above it is easy to work out any other values you may need.

Detect Key Key	Port 5 Sense Value		Port 5 Scan Value	
	Hex	Decimal	Hex	Decimal
A	£DF		£FE	
B	£7F		£FB	
C	£7F		£FD	
D	£DF		£FD	
E	£7F		£FD	
F	£EF		£FB	
G	£DF		£FB	
H	£EF		£F7	
I	£FB		£EF	
J	£DF		£F7	
K	£EF		£EF	
L	£DF		£EF	
M	£7F		£7F	
N	£BF		£F7	
O	£F7		£EF	
P	£FB		£DF	

Q	£F7	£FE
R	£FB	£FB
S	£EF	£FD
T	£F7	£FB
U	£F7	£F7
V	£BF	£FB
W	£FB	£FD
X	£BF	£FD
Y	£FB	£F7
Z	£7F	£FE
0	£FD	£DF
1	£FE	£FE
2	£FD	£FD
3	£FE	£FD
4	£FD	£FB
5	£FE	£FB
6	£FD	£F7
7	£FE	£F7
8	£FD	£EF
9	£FE	£FE
Cursor up	£FB	£7F (see section 5.2 – 5.3)
Cursor down	£BF	£7F
Cursor right	£EF	£7F
Cursor left	£F7	£7F
Home key	£DF	£7F

Appendix F

7.6

Alternative KBD Routine

Please refer to section 5.5

```

;
; The bkd routine returns after finding the first valid character
; If control is pressed then upper and lower case is ignored
; If alpha lock is found then alpha flag is set
; If shift is found then shifted characters are generated
; Numeric keypad is operated by alpha lock
; Numeric keypad produces code depending on NK flag
;
KBFLAG      *AL*TAB/LF*FN/NU*AUTO ON*AUTO SPEED*ICC*
;BIT        7   6   5   4           3 2 1   0
;
; 7 alpha lock on/off
; 6 tab or lf key
; 5 page mode or scroll

```

```

;4 auto repeat on
;3 speed of auto
;0 international codes on alpha lock
;
SENSE1 EQU 5
SENSE 2 EQU 6
DR EQU 5
;
;
KBFLAG: DB 40H
LASTKY: DB 0
;
KBD: PUSH BC
      PUSH DE
      PUSH HL
      CALL SSS ;TEST ROUTINE
      AND A
      DB 0,0,0,0,0,0,0,0,0,0
EXIT: POP HL
      POP DE
      POP BC
      RET
DEB: CALL DEB1
      JR Z,OKI
      LD (LASTKY),A
      CP 128
      RET C
OKI: LD A,0
      LD A, (LASTKY)
      CP B
      LD A,B
      RET
SSS: CALL KBDSTA
      CALL DEB
      RET
KBDSTA: LD A,251
        OUT (DR),A
        IN A, (DR)
        BIT 0,A
        JR JZ,NORMAL
CNTRL: CALL NORMAL
        LD B,A
        CP 128
        JR Z,CN3
        CP 136
        JR NZ,CN2
CN3: CALL DEB1
      RET Z
      LD C,01
      CALL SWITCH
      LD A,B
      RET

```

```

CN2:      CP 129
          JR Z,CN4
          CP 137
          JR NZ,CN1
CN4:      CALL DEB1
          RET Z
          LD C,40H
          CALL SWITCH
          LD A,B
          RET
CN1:      LD A,B
          BIT 6,A
          JR Z,NOCONT
          BIT 7,A
          JR NZ,NOCONT
          AND 1FH
          RET
NOCONT:   LD, A,O
          RET
SWITCH:   LED A, (KBFLAG) ;GET FLAGS
          LD D,A           ;SAVE FLAGS IN D
          AND C           ;A=0 IF BIT NOT SET
          JR Z,BITZ
          CPL             ;BIT IS SET TO UNSET IT
          AND D
          JR SWEND
BITZ:     OR D
          OR C
SWEND:    LD (KBFLAG),A
          RET
NORMAL:   CALL SHIF
          JR Z,NSHIFT
          CALL SCAN
          LD D,A
          LD BC,BASE
          CALL KBDLUK
          LD B,A
          LD A,9KBFLAG0
          LD C,A
          BIT 7,A
          LD A,B
          JR Z,NCONT
AUTO:    ; CP 32
        ; JR C,NSH1
NOTNP:   CP 64
          RET C
          JR NZ,AU96
          BIT 0,C
          RET Z
          ADD A,32
          RET

```

```

AUG96:    CP 95
           JR NC,NSH1
           BIT 0,C
           RET Z
           ADD A,32
           RET
NSH1:     LD A,D           ;PAD ON SHIFT
           JR NSH2
NSHIFT:   CALL SCAN
NSH2:     LD BC,UPPER
           CALL KBDLUK
NCONT:    LD B,A
           CP 144
           JR NZ,NORM1
           LD A,(KBFLAG)
           LAD C,A
           LD A,9
           BIT 6,C
           RET Z
           INC A
           RET
NORM1:    CP 145
           JR NZ,NORM2
           CALL DEB1
           RET \
           LD C,128
           CALL SWITCH
           LD A,145
           RET
NORM2:    CP 28           ;28-SCROLL/PAGE
           RET NZ         ;NORMAL CHAR
           LD 3,A         ;save 28
           CALL DEB1
           RET Z
           LD A,29
           LD A,(KBFLAG)
           BIT 5,A
           JR Z,NORM3
           INC E
NORM3:    LD C,32
           CALL SWITCH
NORM4:    LD A,E
           RET
;MATRIX SCAN
SCAN:     LD B,B         ;B=8=DRIVE COUNTER
           LD C,0         ;SENSE COUNTER
           LD,A,OFFH
           AND A
SCAN2:    RL A           ;A=FE,FD,FB,F7,EF,DF,BF,7F
           PUSH AF        ;SAVE DRIVE
           OUT (DR),A
           IN A,(SENSE1)

```

```

                CP OFFH                ;IF FF THEN NO SENSE, TRY OTHER DRIVE
                JR Z,SCAN3
CHECK1:        PUSH AF                ;SAVE SENSE
CH10:         LD A,2                  ;CHECK FOR SHIFT
                CP B                  ;IF B=2 THEN RESET SHIFTS
                JR NZ,CH11
                POP AF
                SET 6,A
                JR CH13
CH11:         LD A,6                  ;CHECK FOR CNTRL
                CP B
                JR Z,CH12
                POP AF
                JR VALID
CH12:         POP AF                ;RESTORE SENSE
CH13:         SET 0,A                ;RESET ALPHA SHIFT AND CONTROL
CH14:         CP OFFH                ;TRY AGAIN AFTER ELIMINATING ODD KEYS
                JR Z,SCAN3            ;IF STILL 0 THEN TRY OTHER DRIVE
;VALID KEY FOUND
VALID:        POP DE                ;REMOVE DRIVE
                LD C,0
CH15:         RRCA
                JR NC,ENDSCAN        ;B=1 -8, C=0-7
                INC C
                JR CH15
SCAN3         IN A, (SENSE 2)        ;SECOND DRIVE
                AND 03
                CP 3
                JR Z,SCAN4
CHECK2:       AND A
                ADD A,7
                LD C,A                ;C=8 OR 9
                POP AF
                JR ENDSCAN
SCAN4:        POP AF                ;RESTORE DRIVE
                DJNZ SCAN2            ;MOVE TO NEXT DRIVE
                LD C,0
ENDSCAN:     LD A,C
                SLA A
                SLA A
                SLA A
                ADD A,B
                AND A
                RET
SHIF:        LD A,191
                OUT (DR), A          ;Z IMPLIES SHIFT
                IN A,(DR)
                BIT 6,A
                RET Z
                BIT 0,A
                RET

```

```

KBDLUK:   LD H,0
          LD L,A
          ADD HL,BC
          LD A,(HL)
          RET
;
BASE:     DB 00
DB 122,0,97,145,113,0,27,49
DB 99,120,100,115,101,119,50,51
DB 98,118,103, 102,116,114,52,53
DB 109,110, 106,104,117,131,54,55
DB 46, 44,108,107,111,105,56,57
DB 95,47,58,59,64,112,48,45
DB 14,0,13,93,144,91,94,92
DB 12,10,26,25,8,11,5,28
DB 134,135,132,133,131,130,129,128
DB 32,0,0,0,17,9,8,3
UPPER:    DB 0
DB 90,0,65,145,81 0,27,33
DB 67,88,68,83,69,87,34,35
DB 66,86,71,70,84,82,36,37
DB 77,78,74,72,85,89,38,39
DB 62,60,76,75,79,73,40,41
DB 95,63,42,43,96,80,48,61
DB 48,0,13,125,144,123,126,124
DB 13,46,50,51,49,53,56,55
DB 142,143,140,141,139,138,137,136
DB 32,0,0,0,54,52,8,57

```

Appendix G

7.7

Alternative INOUT Routine

Please refer to section 6.2

This appendix is a straight forward listing of an edited version of the source code used to save/load data bytes.

it is intended for Z80 machine code programmers only.

The same process can be produced by using an absolute call and the parameter set up explained in section 6.

```

;
CASSETTE ROUTINES
;
CASSET:   JP INOUT           ;ENTRY POINT -- (TYPE) = 0 for saving
;                                           = 1 for loading
;
;                                           ;HL = start address,  DE = no of bytes
;

```

```

CASPORT      EQU      3+OFFSET
SNDPTI      EQU      3+OFFSET
DRIVE       EQU      5+OFFSET
SNDPT0      EQU      6+OFFSET
SENSE2      EQU      6+OFFSET
PORT        EQU      8+OFFSET
CASonOF     EQU      1FH
DELAY       EQU      1500
;
;
;***** VARIABLES THAT YOU WILL NEED *****
;
;
CASBAUD:    DB 40H
MIDVAL:     DB 0B0H
TYPE:       DS 1
;
;IJHIGH/IJLOW POINTERS TO START OF CASSET ROUTINE
;(IE THE I/O POINT CASSET)
;USED BY CTC CHIP AND MUST BE ON AN 8 BYTE BOUNDARY
;
;          *****
;
IJHIGH      EQU      (HIGH BYTE)
IJLOW      EQU      (LOW BYTE)
;
;*****
;
;*****START OF CASSETTE INTERFACE*****
;
;
;Interrupt routine – toggles carry flag
;
TOGGLE:     CCF
            EI
            RETI
;
;
BLIP:
            EX AF, AF'
            XOR OFH
            OUT (SNDPTO),A ;makes a noise
            EX AF,AF'
            IN A,(SNDPTI)
            RET
;
;
;
INBIT::
            CALL BLIP
            EI
            XOR A          ;Carry flag cleared

```

```

INBIT1:      JR NC,INBIT 1
INBIT2:      DEC A
              JP C,INBIT2      ;A bit faster
              DI
              CP (IX+0)        ;IX = MIDVAL
              RET

;
;
;
;
OUTBIT::     CALL BLIP          ;Preserves carry flag
              LD A,0           ;Preserves carry flag
              JR C,HIGH

;
;
;
;
;Both LOW and HIGH entered with A = 0
;
LOW::        JR NC,LOW
              OUT (CASPORT),A
LOW1:        JR C,LOW1
              INC A
              OUT (CASPORT),A
              RET

;
;
;
;
HIGH::       JR C,HIGH
HIGH1:       JR NC,HIGH1
              OUT (CASPORT),A
HIGH2:       JR C,HIGH2
HIGH3:       JR NC,HIGH3
              INC A
              OUT (CASPORT),A
              RET

;
;
;
;
INBYTE::     LD B,B
INBY1:       CALL INBIT
              RR C
              DJNZ INBY1
              RET

;
;
;
;
OUTBYTE::    LD B,B
OUTBY1:      RR C
              CALL OUTBIT
              DJNZ OUTBY1
              RET

;

```

```

;
;
;
INOUT::      LD A,D
              OR E
              RET Z          ;Return if size of block = 0
              CALL SETint
              EX AF,AF'
              LD A,128+16

EX AF,AF
              JR Z,OUTBLOCK  ;Jump if saving

;
;
;
INBLOCK::    LD B,O
INBLK1:      CALL INBIT
              JR C,INBLOCK
              DJNZ INBLK1    ;Look for 50 low bits

STBIT:      EI
              XOR A
              CCF           ;Set carry
              CALL INBIT2   ;Find long pulse
              JR NCSTBIT

INBLK2::     CALL INBYTE    ;C = byte from tape
              LD (HL),C     ;Store byte
              INC HL
              DEC DE
              LD A,D
              OR E
              JR NZ,INBLK2
              JR RESint

;
;
;
;
OUTBLOCK::   LD BC, DELAY
OTBLK1:      XOR A
              CALL LOW
              DEC BE
              LD A,B
              OR C
              JR NZ,OTBLK11  ;Carry flag reset A = 0
DEL1:        JR NC,DEL1     ;Short high
              OUT (CASPORT),A ;A = 0
DEL2:        JR C,DEL2     ;One more interrupt
              CCF
              CALL HIGH2

OTBLK2:      LD C,(HL)     ;C = byte to be output
              CALL OUTBYTE
              INC HL
              DEC DE
              LD A,D
              OR E
              JR NZ,OTBLK2

```

```

;
;
;
RESint:          CALL INJINIT
                 LD A,55H
                 OUT (CASonOFF),a
                 RET

;
;
;
;Routine to set CTC interrupts going
;Routine Z if saving, NZ if loading/verifying
;
SETint:          CALL IJINIT
                 PUSH HL
                 LD HL,TOGGLE
                 LD (IJTABLE+2),HL
                 LD (IJTABLE+6),HL
                 POP HL
                 LD IX,MIDVAL
                 LD A,0AAH
                 OUT (CASonOFF),A
                 LD C,PORT+3
                 LD B,0C5H           ;Assume loading
                 LD A,(TYPE)
                 AND A
                 JR NZ,SETin1       ;Jump if A = 1 ie loading
                 LD C,PORT+1
                 LD A, (CASBAUD)
SETin1:          OUT (C),B
                 OUT (C),A
                 EI
                 RETI               ;Clear interrupts and return

;
;Init CTC and set all channels to
;disable interrupts and software reset
;
ijinit::         DI
                 IM      2
                 LD      A,IJHIGH
                 LD      I,A
                 LD      A,IJLOW
                 OUT     (PORT),A
KILLINT::        LD      A,3
KILLO:          OUT     (PORT),A
                 OUT     (PORT+1),A
                 OUT     (PORT+2),A
                 OUT     (PORT+3),A
                 RETI

```

```
;
;
;*****END OF CASSETTE INTERFACE *****
;
;          DB    0,0          ;DUMMY PAD OUT BYTES
;
INJTABLE:: DB    "IJTABLE"
```